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A SYSTEM OF EXTENDED SET THEORY

by

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## PREFACE

A system  $K$  of set theory is defined which contains classes, called "multiplicities," and sets, called "extended classes." Certain extended classes, called "extended sets," are singled out; these are characterized by the property that their cardinalities are less than a fixed cardinal number  $\aleph_0$ , which is shown to be strongly inaccessible. The terminology and the axioms are those of Friedman, J.I.: Proper Classes as Members of Extended Sets, Math. Ann. 183, 232-240 (1969), with the maximality axiom replaced by the cardinality restriction on extended sets. Analogously to the von Neumann-Bernays-Gödel system (NBG), the system  $K$  is obtainable from an extended version of the von Neumann function. In fact, for each ordinal number  $\beta$ , the extended classes  $M_\beta$  and  $K_\beta$  given by the axiom of universes, and consisting respectively of extended sets and proper extended classes, coincide respectively with extended classes  $S_\beta$  and  $C_\beta$  arising from the inductive definition of the extended von Neumann function. Some results on cardinal numbers are obtained which involve the  $\aleph_\beta$ ,  $M_\beta$  and  $K_\beta$ , where the  $\aleph_\beta$  are the ordinal numbers of  $K$  which are not extended sets. Again analogously with NBG, the extended von Neumann function together with the inaccessible ordinals  $\aleph_\beta$  with  $\beta > 0$  (if any) yield the natural models of  $K$ . The system  $K$  is relatively consistent with NBG together with the hypothesis of the existence of two inaccessible ordinals. Gödel's construction which shows the relative consistency of the axiom of constructibility with NBG is extended to  $K$ .

Bracketed numbers refer to the bibliography. Within each section a number cited in a proof refers to an earlier theorem in that section. All notation used may be found in [1], [2] and [3].

I wish to thank Professor Elliott Mendelson of Queens College of the City University of New York for suggesting that I axiomatize and study

a system based upon the extended von Neumann function, and for suggesting much of the reading I have done in set theory and mathematical logic over several years.

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1. Introduction. The first objective is to modify the extended set theory of [1] so that any extended set of the resulting theory K will have the same cardinality as some extended set of  $M_0$ . by analogy with the von Neumann function of NBG, the extended von Neumann function  $\psi$  given below (without formal justification) is required to be definable as a function of K, and to yield all of the extended classes and extended sets of K. The ordinal numbers  $\alpha$  in the definition of  $\psi_\alpha$  may be regarded as those of  $M_0$ ; note that  $\psi_0$  is the ordinary von Neumann function.

$$\begin{aligned} \psi_0' 0 &= 0 & \psi_\beta' 0 &= \bigcup_{\gamma < \beta} S_\gamma \cup C_\gamma \\ \psi_0' \alpha + 1 &= P(\psi_0' \alpha) & \psi_\beta' \alpha + 1 &= \{ \mu : \mu \leq \psi_\beta' \alpha \wedge \mu < S_0 \} \\ \text{lim}(\lambda) \rightarrow \psi_0' \lambda &= \bigcup \psi_0'' \lambda & \text{lim}(\lambda) \rightarrow \psi_\beta' \lambda &= \bigcup \psi_\beta'' \lambda \\ S_0 &= \bigcup \psi_0'' \theta_{\sim M} & S_\beta &= (\bigcup \psi_\beta'' \theta_{\sim M}) - \bigcup_{\gamma < \beta} C_\gamma \\ C_0 &= P(\bigcup \psi_0'' \theta_{\sim M}) - S_0 & C_\beta &= P(\bigcup \psi_\beta'' \theta_{\sim M}) - S_\beta \end{aligned}$$

Write  $U_\alpha$  for  $\bigcup \psi_\alpha'' \theta_{\sim M}$ ,  $P_\alpha$  for  $P(\bigcup \psi_\alpha'' \theta_{\sim M})$ . Then by induction,

$$\begin{aligned} S_\alpha &= U_\alpha - \bigcup_{\lambda < \alpha} (P_\lambda - U_\lambda), \\ C_\alpha &= P_\alpha - S_\alpha = P_\alpha - (U_\alpha - \bigcup_{\lambda < \alpha} (P_\lambda - U_\lambda)), \\ \therefore S_\alpha \cup C_\alpha &= (U_\alpha - \bigcup_{\lambda < \alpha} (P_\lambda - U_\lambda)) \cup P_\alpha, \\ \therefore \bigcup_{\alpha < \beta} S_\alpha \cup C_\alpha &= \bigcup_{\alpha < \beta} (U_\alpha - \bigcup_{\lambda < \alpha} (P_\lambda - U_\lambda)) \cup \bigcup_{\alpha < \beta} P_\alpha = \bigcup_{\alpha < \beta} U_\alpha \cup \bigcup_{\alpha < \beta} P_\alpha, \text{ i.e.,} \\ \bigcup_{\alpha < \beta} S_\alpha \cup C_\alpha &= \bigcup_{\alpha < \beta} (\bigcup \psi_\alpha'' \theta_{\sim M}) \cup \bigcup_{\alpha < \beta} P(\bigcup \psi_\alpha'' \theta_{\sim M}). \end{aligned}$$

2. Definition of the theory K. K is a first-order theory with a 2-place predicate  $\epsilon$  and  $\hat{1}$ -place predicate M. Upper-case letters are used for the individual variables and will be said to denote "multiplicities." Lower-case letters are used for individual variables which denote multiplicities X such that  $(\exists Y)(X \epsilon Y)$ . Such multiplicities will be called "extended classes;" their properties are given by the "general axioms," which in form are the axioms for NBG. (Thus the multiplicities of K correspond to the classes of NBG, and the extended classes of K to the sets of NBG.) The comprehension schema of NBG gives

the existence of the multiplicity  $\{x:MX\}$ , also denoted by  $M$ , and the multiplicity  $\{x:x=x\}$ , denoted by  $V$ . If  $x \in M$ ,  $x$  is called an "extended set." The extended sets satisfy the "special axioms" (in addition to the "general axioms"); the first six of these are as in [1].

General axioms:

Extensionality  $X=Y \rightarrow (X \in Z \leftrightarrow Y \in Z)$ .

Pairing  $(x)(y)(\exists z)(u)(u \in z \leftrightarrow u=x \vee u=y)$ .

Class existence axioms B1-B7 as in [2].

Sum  $(x)(\exists y)(u)(u \in y \leftrightarrow (\exists v)(u \in v \wedge v \in x))$ .

Replacement  $(x)(\forall \eta(X) \rightarrow (\exists y)(u \in y \leftrightarrow (\exists v)(\langle v, u \rangle \in X \wedge v \in x)))$ .

Power  $(x)(\exists y)(u)(u \in y \leftrightarrow u \subseteq x)$ .

Strong axiom of choice  $(\exists X)(\text{Fnc}(X) \wedge (u)(u \neq \emptyset \rightarrow X' u \in u))$ .

Restriction  $(X)(X \neq \emptyset \rightarrow (\exists Y)(Y \in X \wedge Y \cdot X = \emptyset))$ .

Infinity  $(\exists x)(\emptyset \in x \wedge (u)(u \in x \rightarrow u \cup \{u\} \in x))$ .

Special axioms:

Subsets<sub>M</sub>  $x \in M \rightarrow x \cdot Z \in M$ .

Union<sub>M</sub>  $x \in M \rightarrow \cup_M x \in M$  ( $\cup_M x = \cup(x \cdot M)$ ).

Power<sub>M</sub>  $x \in M \rightarrow P_x \in M$ .

Pairing<sub>M</sub>  $x, y \in M \rightarrow \{x, y\} \in M$ .

Infinity<sub>M</sub>  $(\exists x)(x \in M \wedge \exists x \in M \cdot \emptyset \in x \wedge (u)(u \in x \rightarrow u \cup \{u\} \in x))$  ( $\exists x(x = \cup x \cup \cup \cup x \cup \dots)$ ).

Universes  $\{x:Tx \subseteq M \cup z\} \in V$ .

Cardinality  $x \in M \leftrightarrow \exists y \in M, Ty \subseteq M \wedge x \approx y$ . (This is the requirement that extended sets have "small" cardinality.)

Notation: The axiom of universes permits the definition of the extended classes  $K_\alpha = \{x \in V - M : Tx \subseteq M \cup \cup_{\gamma < \alpha} K_\gamma\}$  and  $M_\alpha = \{x \in M : Tx \subseteq M \cup \cup_{\gamma < \alpha} M_\gamma\}$ .

3. Some immediate results. As in [1], the following are proved from the general axioms and the first six of the special axioms.

1.  $K_\gamma \cdot M_\alpha = \emptyset$ ;  $i \leq \beta \rightarrow (K_i \subseteq K_\beta ; M_i \subseteq M_\beta)$ .

2.  $V = \cup_{\alpha \in \theta_\alpha} M_\alpha \cup K_\alpha$ ;  $M = \cup_{\alpha \in \theta_\alpha} M_\alpha$ ;  $V - M = \cup_{\alpha \in \theta_\alpha} K_\alpha$ .

3.  $M_A \in K_A; M_B \in K_A$  (proof uses axiom of choice).

4.  $x, y \in M_A \rightarrow x \cdot z, U_n x, P x, \{x, y\} \in M_A$ .

4. Formal definition of the extended von Neumann function.

1. Let  $On_M = On \cdot M, On_{M_0} = On \cdot M_0$ . Then  $On_M = On_{M_0}$ . (For otherwise, E least  $x \in On_M - On_{M_0}$ . Then  $x \in M \rightarrow \exists w \in M_0, x = w$ .  $\therefore y < x \rightarrow \exists v \in w, y < v$ . Since  $Tv \in Tw, v \in M_0$ .  $\therefore y \in M_0$ .  $\therefore x \in On_{M_0} \subseteq M_0 \subseteq M$ . If  $z \in Ux, \exists y \in x, z \in y$ .  $\therefore z \in x$ , since  $z < y < x$ .  $\therefore Ux \subseteq x \subseteq M$ . In general,  $Tx \subseteq M$ . Since  $x \in M, x \in M_0, \rightarrow \dots$ )  $\therefore On_M = On_{M_0} = On \cdot M_0 \subseteq M_0 \subseteq V$ , so  $On_M \in V$ .

2. There is a function  $\psi$  with domain  $On_M, \psi' 0 = 0, \psi' \alpha + 1 = P(\psi' \alpha), \text{Lim}(\alpha) \rightarrow \psi' \alpha = U \psi'' \alpha$ . (Proof is by the general axioms, i.e., by NHG. By 1.,  $On_M$  is an extended class, and by definition of  $On$  and  $M_0, u \in On_{M_0} \wedge w \in u \rightarrow w \in On_{M_0}$ .  $\therefore On_{M_0}$  is an ordinal number.) Let  $S_0 = U \psi'' On_M$ .

3.  $(X_0)(X_1)(X_2)(E, Y)((\exists \beta \in Y) \cup (Y = On \times On_M \wedge (\beta)(Y' < \beta, 0) = X_0'(\beta \times On_M \upharpoonright Y) \wedge (=)(Y' < \beta, \alpha') = X_1' Y' < \beta, \alpha') \wedge (\text{Lim}(\alpha) \rightarrow Y' < \beta, \alpha') = X_2'(\{\beta\} \times \alpha \upharpoonright Y)))$ . (Proof is by the general axioms; the fact  $On_M \in V$  is used.)

Proof: Let  $Y_1 = \{u; \text{Func}(u) \wedge (E v \in On)(D(u) = v \times On_M \wedge (\beta)(\beta < v \rightarrow u' < \beta, 0) = X_0'(\beta \times On_M \upharpoonright u) \wedge (\alpha \in On_M)(u' < \beta, \alpha') = X_1' u' < \beta, \alpha') \wedge (\text{Lim}(\alpha) \rightarrow u' < \beta, \alpha') = X_2'(\{\beta\} \times \alpha \upharpoonright u))\}$ . If  $u_1, u_2 \in Y_1$ , then either  $u_1 \subseteq u_2$  or  $u_2 \subseteq u_1$ . For let  $v_1 = On_M \times D(u_1), v_2 = On_M \times D(u_2)$ . Assume  $v_1 \subseteq v_2$ . Let  $w = \{\alpha; \alpha < v_1 \wedge (E \delta \in On_M)(u_1' < \alpha, \delta) \neq u_2' < \alpha, \delta)\}$ . If  $w \neq \emptyset$  let  $\kappa$  be the least ordinal in  $w$ . Then  $u_1' < \kappa, 0) = X_0'(\kappa \times On_M \upharpoonright u_1) = X_0'(\kappa \times On_M \upharpoonright u_2) = u_2' < \kappa, 0)$ . If  $\exists \alpha, u_1' < \kappa, \alpha) \neq u_2' < \kappa, \alpha)$ , let  $\epsilon$  be the least such. Then  $\epsilon > 0$ . If  $\epsilon = \nu'$ ,  $u_1' < \kappa, \nu) = X_1' u_1' < \kappa, \nu) = X_1' u_2' < \epsilon, \nu) = u_2' < \kappa, \epsilon)$ . If  $\text{Lim}(\epsilon)$ , then  $u_1' < \kappa, \epsilon) = X_2'(\{\kappa\} \times \epsilon \upharpoonright u_1) = X_2'(\{\kappa\} \times \epsilon \upharpoonright u_2) = u_2' < \kappa, \epsilon)$ . Hence  $(\delta \in On_M)(u_1' < \kappa, \delta) = u_2' < \kappa, \delta)$ .  $\therefore u_1 = (v_1 \times On_M) \upharpoonright u_1 = (v_1 \times On_M) \upharpoonright u_2 \subseteq u_2$ . Thus any two functions in  $Y_1$  agree on their common domain. Let  $Y = \Pi Y_1$ . ( $Y$  is a multiplicity as it is defined by a predicative formula.) Then  $Y$  is a function. For  $Y \subseteq V \times V$ ; and  $U \cap Y$ , since  $\langle \alpha, \beta \rangle \in Y \wedge \langle \alpha, \beta \rangle \in Y \rightarrow \exists u_1, u_2 \in Y, \langle \alpha, \beta \rangle \in u_1 \wedge \langle \alpha, \beta \rangle \in u_2$ ; but  $u_1 \subseteq u_2$  or  $u_2 \subseteq u_1$ . For some  $\beta \in On, D(Y) = \beta \times On_M$ , or  $D(Y) = On \times On_M$ . For  $D(Y) = Z \times On_M$  for some  $Z \subseteq On$ . Let  $x \in Z$ . Show  $x = \text{Seg}(Z, x)$ .

Since  $x \in On$ ,  $\{w: w \in Z \wedge w < x\} \subseteq x$ . Let  $w \in x$ .  $\exists u \in Y, \alpha \in On_\mu, \langle x, \alpha \rangle \in D(u)$ .  $\therefore \langle w, \alpha \rangle \in D(u)$ .  $\therefore w \in Z$ .  $\therefore x \text{ Seg}(Z, x)$ .  $\therefore Z \in On \vee Z = On$ . Suppose  $Z \in On$ . Let  $W = Y \cup \{\langle \langle Z, 0 \rangle, X_0 \uparrow (Z \times On_\mu \uparrow Y) \rangle\} \cup \{\langle \langle Z, \alpha' \rangle, X_1 \uparrow Y \uparrow \langle Z, \alpha \rangle \rangle : \alpha \in On_\mu\} \cup \{\langle \langle Z, \alpha \rangle, X_2 \uparrow (\{Z\} \times \alpha \uparrow Y) \rangle : \alpha \in On_\mu \wedge \text{Lim}(\alpha)\}$ . Then  $W \in Y$ , so  $D(W) \subseteq D(Y)$ . But  $D(Y) = Z \times On_\mu \subseteq Z' \times On_\mu = D(W)$ ,  $\rightarrow \leftarrow$ .

Uniqueness of Y: Suppose Z also has the defining property of Y. If  $Y \neq Z$ , let  $\alpha$  be the first ordinal,  $\exists \beta \in On, Y \uparrow \langle \beta, \alpha \rangle \neq Z \uparrow \langle \beta, \alpha \rangle$ . Then  $Y \uparrow \langle \beta, 0 \rangle = X_0 \uparrow (\beta \times On_\mu \uparrow Y) = X_0 \uparrow (\beta \times On_\mu \uparrow Z) = Z \uparrow \langle \beta, 0 \rangle$ .  $\therefore \alpha \neq 0$ . If  $\alpha = \gamma'$ ,  $Y \uparrow \langle \beta, \alpha \rangle = X_1 \uparrow Y \uparrow \langle \beta, \gamma' \rangle = X_1 \uparrow Z \uparrow \langle \beta, \gamma' \rangle = Z \uparrow \langle \beta, \alpha \rangle$ .  $\therefore \alpha \neq \gamma'$ . If  $\text{Lim}(\alpha)$ ,  $Y \uparrow \langle \beta, \alpha \rangle = X_2 \uparrow (\{\beta\} \times \alpha \uparrow Y) = X_2 \uparrow (\{\beta\} \times \alpha \uparrow Z) = Z \uparrow \langle \beta, \alpha \rangle$ .  $\therefore \sim \text{Lim}(\alpha)$ .  $\therefore Y = Z$ .

4. In 3., take  $X_0 = \{\langle u, v \rangle : v = U \cap R(u) \cup \bigcup_{x \in D(u)} P(U \cap R(u \cdot (\{x\} \times V \times V)))\}$ ,  $X_1 = \{\langle u, v \rangle : v = \{x : x \leq u \wedge x \in S_0\}\}$ ,  $X_2 = \{\langle u, v \rangle : v = U \cap R(u)\}$ . Then Y is the required function  $\varphi$ . The notation  $\varphi' \alpha$  is used for  $\varphi \uparrow \langle \beta, \alpha \rangle$ . By NBG,  $u \in S_0 \rightarrow u \in S_0$ , so that  $\varphi_0 = \psi$ . ( $u \in S_0 \rightarrow \exists$  minimal  $\alpha$ ,  $u \in \psi' \alpha$ . If  $\psi' \alpha \neq S_0$ , then  $F(\psi' \alpha) \supset S_0$ . But  $F(\psi' \alpha) = \psi' \alpha + 1 \subseteq S_0$ ,  $\rightarrow \leftarrow$ .)

5. From the general axioms it follows that the extended classes  $S_\alpha \cup C_\alpha$  constitute a multiplicity, and hence there is a multiplicity  $H = \bigcup_{\alpha \in On} S_\alpha \cup C_\alpha$ .

5. Some properties of the  $S_\alpha$  and  $C_\alpha$ ;  $S_\alpha = M_\alpha$  and  $C_\alpha = K_\alpha$ .

1.  $\alpha < \beta \rightarrow \bigcup \varphi_\alpha \text{ " } On_\mu \subseteq \bigcup \varphi_\beta \text{ " } On_\mu$ .

Proof: Let  $x \in \bigcup \varphi_\alpha \text{ " } On_\mu$ . If  $x \in S_\alpha$  then  $x \in \varphi_\beta \uparrow 0 \subseteq \bigcup \varphi_\beta \text{ " } On_\mu$ . If  $x \notin S_\alpha$ , then by definition of  $S_\alpha$ ,  $x \in C_\gamma$  for some  $\gamma < \alpha$ .  $\therefore x \in \varphi_\beta \uparrow 0 \subseteq \bigcup \varphi_\beta \text{ " } On_\mu$ .

2.  $(\alpha)(\beta)(C_\alpha \cdot S_\beta = \emptyset)$ .

Proof: If  $\alpha = \beta$ ,  $C_\alpha \cdot S_\beta = \emptyset$  by definition of  $C_\alpha$ . If  $\alpha < \beta$ ,  $C_\alpha \cdot S_\beta = \emptyset$  by definition of  $S_\beta$ . Assume  $\beta < \alpha$ . True if  $\alpha = 0$ . Assume for ordinals  $< \alpha$ . Suppose  $x \in S_\beta \cdot C_\alpha$ . Then  $x \in C_\alpha \rightarrow x \notin S_\alpha$ . By 1.,  $x \in S_\beta \rightarrow x \in \bigcup \varphi_\alpha \text{ " } On_\mu$ . By definition of  $S_\alpha$ ,  $x \notin S_\alpha \rightarrow x \in C_\gamma$  for some  $\gamma < \alpha$ . If  $\gamma \leq \beta$  then  $x \in S_\beta \cdot C_\gamma = \emptyset$  by the previous cases,  $\rightarrow \leftarrow$ . If  $\beta < \gamma$ ,  $x \in S_\beta \cdot C_\gamma = \emptyset$  by induction hypothesis,  $\rightarrow \leftarrow$ .

3.  $\alpha < \beta \rightarrow S_\alpha \subseteq S_\beta \wedge C_\alpha \subseteq C_\beta$ .

Proof: If  $x \in S_\alpha$  then  $x \in \varphi_\beta \uparrow 0 \subseteq \bigcup \varphi_\beta \text{ " } On_\mu$ , and  $x \notin \bigcup_{\gamma < \beta} C_\gamma$  by 2.  $\therefore x \in S_\beta$ . If  $x \in C_\alpha$

then  $x \in \cup_{\beta \in \sigma} \psi_\beta \text{On}_\mu \cup \psi_\beta \text{On}_\mu$ , and  $x \notin S_\beta$  by 4.  $\therefore x \in C_\beta$ .

4.  $u \in H \rightarrow u \in H$ .

Proof: Let  $u \in V \wedge u \in \cup_{\beta \in \sigma} S_\beta \cup C_\beta$ . If  $x \in u$ ,  $\exists$  least  $\beta$ ,  $x \in S_\beta \cup C_\beta$ . By the axiom of replacement, these  $\beta$  form an extended class. Let  $\lambda$  be greater than all such  $\beta$ . Then  $x \in u \rightarrow x \in S_\lambda \cup C_\lambda \subseteq S_\lambda \cup C_\lambda$ , by 3.  $\therefore u \in S_\lambda \cup C_\lambda = \psi_{\lambda+1} \cup 0$ . If  $u \notin S_{\lambda+1}$ , then  $u \in P(\cup_{\beta \in \sigma} \psi_{\beta+1} \text{On}_\mu) - S_{\lambda+1} = C_{\lambda+1}$ .  $\therefore u \in \cup_{\beta \in \sigma} S_\beta \cup C_\beta = H$ .

5.  $V = H$ .

Proof: If  $V \neq H \neq \emptyset$ ,  $\exists y \in V - H, y \cdot (V - H) = \emptyset$ . Hence by 4.,  $y \in H$ .

6.  $S_0 \subseteq M$ .

Proof: (i) Assertion:  $\beta \in \text{On}_\mu \rightarrow \psi_0 \beta \in M$ . (Proof by induction.)

(a)  $\beta = 0$ . In  $\text{Subsets}_\mu$  take  $x \in M$ ,  $x$  is infinite and  $z = 0$ .

(b)  $\beta = \gamma + 1$ . By  $\text{power}_\mu$ ,  $\psi_0 \beta = P(\psi_0 \gamma) \in M$ .

(c)  $\text{Lim}(\beta)$ . Since  $\beta \in V$ ,  $\{\psi_0 \delta : \delta < \beta\} \in M$  by Cardinality. Also,  $\{\psi_0 \delta : \delta \in \beta\} \in M$  by induction hypothesis.  $\therefore \psi_0 \beta \in M$  by Union $_\mu$ .

(ii) Assertion:  $\beta \in \text{On}_\mu \rightarrow \psi_0 \beta \in M$ . (Proof by induction.)

(a)  $\beta = 0$ . True  $\therefore \psi_0 0 = 0$ .

(b)  $\beta = \gamma + 1$ . If  $x \in \psi_0 \beta$  then  $x \subseteq \psi_0 \gamma$ . By (i),  $\psi_0 \gamma \in M$ , so  $x \in M$  by  $\text{Subsets}_\mu$ .

(c)  $\text{Lim}(\beta)$ . Let  $x \in \psi_0 \beta$ .  $\exists \gamma < \beta, x \subseteq \psi_0 \gamma$ . By induction hypothesis,  $\psi_0 \gamma \in M$ , so  $x \in M$ .

7.  $S_0 \cong \text{On}_\mu$ .

Proof: By induction on  $\alpha$ ,  $\alpha \in \text{On}_\mu \rightarrow \alpha \subseteq \psi_0 \alpha + 1$ .  $\therefore \text{On}_\mu \subseteq S_0$ . But  $S_0 \subseteq \text{On}_\mu \cong \text{On}_\mu$ .

For let  $x \in S_0$ .  $\exists$  minimal  $\alpha$ ,  $x \in \psi_0 \alpha$ . By 6.(i),  $\psi_0 \alpha \in M$ , so  $\exists \gamma$  which well-orders  $\psi_0 \alpha$ . By Cardinality,  $\gamma \cong \psi_0 \alpha \rightarrow \gamma \in M$ . Let  $x$  correspond to  $\langle \alpha, \beta \rangle$ , where  $\beta$  is associated with  $x$  by the well-ordering of  $\psi_0 \alpha$ . Thus  $S_0 \cong \text{On}_\mu \times \text{On}_\mu \cong \text{On}_\mu$ .  $\therefore S_0 \cong \text{On}_\mu$ .

8.  $u \in \cup_{\gamma \in \sigma} C_\gamma$  iff  $u \in S_0$ . (Hence  $u \in \cup_{\gamma \in \sigma} S_\gamma$  iff  $u \in S_0$ , by 5.)

Proof: Let  $u \in \cup_{\gamma \in \sigma} C_\gamma$ .  $\exists$  minimal  $\gamma$ ,  $u \in C_\gamma$ . For each  $x \in u$   $\exists$  minimal  $\alpha$ ,  $x \in \psi_0 \alpha$ . Let  $A$  be the extended class of such  $\alpha$ . Then  $A \subseteq u$ . If  $A \cong \text{On}_\mu$ , then  $S_0 \subseteq u$  by 7. Suppose  $A \not\cong \text{On}_\mu$ . Since  $\text{On}_\mu$  is an ordinal number,  $A \cong \delta$

for some  $\lambda \in \text{On}_\mu \approx M_0$ . Since  $A \in M_0$ ,  $T_A \in M_0$ , so  $A \in M$ .  $\cup A \in M$  by Union $_M$ , so that  $A$  has a strict upper bound  $\lambda \in \text{On}_\mu$ . May assume  $\text{Lim}(\lambda)$ . Then  $u \in \varphi_\nu^{-1} \lambda$ . If  $u \in S_0$ , then  $u \in \varphi_\nu^{-1} \lambda + 1$ . Since  $u \in \bigcup_{\tau < \nu} C_\tau$ ,  $u \in S_\tau$ ,  $\dots$ ,  $u \in S_0$ .

Conversely, assume  $u \in \bigcup_{\tau < \nu} S_\tau$ .  $\exists$  minimal  $\gamma$ ,  $u \in S_\gamma$ .  $\exists$  minimal  $\alpha$ ,  $u \in \varphi_\nu^{-1} \alpha + 1$ .  $\therefore u \in \varphi_\nu^{-1} \alpha \cup S_0$ .

9.  $M = \bigcup_{\alpha < \theta} S_\alpha$ . (Thus  $V = M \cup \bigcup_{\alpha < \theta} C_\alpha$ , by 5.)

Proof: (i) Assertion:  $\alpha \in \text{On} \rightarrow S_\alpha \in M$ . True if  $\alpha = 0$ , by c. Assume for all  $\beta < \alpha$ .

Let  $x \in S_\alpha$ . By induction hypothesis, may assume  $\alpha$  minimal,  $x \in S_\alpha$ .  $\exists$  minimal  $\beta$ ,  $x \in \varphi_\alpha^{-1} \beta$ . By minimality of  $\alpha$ ,  $\beta \neq 0$ , and by minimality of  $\beta$ ,  $\sim \text{Lim}(\beta)$ .  $\therefore \beta = \gamma + 1$ .  $\therefore x \in \varphi_\alpha^{-1} \gamma \cup S_0$ .  $\therefore \exists y \in S_0, y \in S_0 \wedge x \in y$ . By 3.,  $y \in C_0 \rightarrow y \in S_0$ .  $\therefore y \in S_0 \in M$ .  $\therefore x \in M$  by Cardinality.

(ii) Assertion:  $M \in \bigcup_{\alpha < \theta} S_\alpha$ . Let  $x \in M$ .  $\exists \alpha$  which well-orders  $x$ , so  $x \in \alpha$ . By Cardinality,  $\alpha \in M$ .  $\therefore \alpha \in \text{On}_\mu \approx S_0$ . Since  $\alpha \in S_0, \alpha \in S_0 \dots x \in S_0 \dots x \in \bigcup_{\alpha < \theta} S_\alpha$ , by 6., i.e.,  $M_0 \subseteq \bigcup_{\alpha < \theta} S_\alpha$ . Let  $y \in M$ . By Cardinality,  $\exists z \in M, Tz \subseteq M \wedge y \in z$ . Since  $z \in M_0, z \in S_0$  by 3.  $\therefore y \in \bigcup_{\alpha < \theta} S_\alpha$  by 5.

10. (a)  $\alpha \in \text{On} \rightarrow C_\alpha \supset S_0$  (hence  $C_\alpha \notin M$ ); (b)  $\beta < \alpha \rightarrow C_\beta \in C_\alpha$ .

Proof: (a)  $C_\alpha \supset C_0 = P(S_0) - S_0$ , and  $P(S_0) \supset S_0$ .

(b)  $C_\beta \subseteq \varphi_\beta^{-1} 0 \in \bigcup \varphi_\alpha^{-1} \text{On}_\mu$ , and  $C_\beta \notin S_\alpha$  by (a).

11.  $\bigcup \varphi_\alpha^{-1} \text{On}_\mu$  is transitive.

Proof: True for  $\alpha = 0$ . Assume for ordinals  $< \alpha$ . Let  $x \in \bigcup \varphi_\alpha^{-1} \text{On}_\mu$ . Then  $x \in \varphi_\alpha^{-1} \tau$  for some minimal  $\tau$ . By definition of  $\varphi$ ,  $\tau$  is not a limit ordinal. If  $\tau = 0$  then  $\exists \beta < \alpha, x \in S_\beta \cup C_\beta$ . If  $x \in S_\beta, x \in \bigcup \varphi_\beta^{-1} \text{On}_\mu$ . By induction hypothesis and 1.,  $x \in \bigcup \varphi_\beta^{-1} \text{On}_\mu \subseteq \bigcup \varphi_\alpha^{-1} \text{On}_\mu$ . If  $x \in C_\beta$ , then  $x \in \bigcup \varphi_\beta^{-1} \text{On}_\mu \subseteq \bigcup \varphi_\alpha^{-1} \text{On}_\mu$ . If  $\tau = n + 1$  and  $x \in \varphi_\alpha^{-1} \tau$ , then  $x \in \varphi_\alpha^{-1} n \cup \bigcup \varphi_\alpha^{-1} \text{On}_\mu$ .

12.  $x \in S_\beta \cup C_\beta \rightarrow T_x \subseteq S_\beta \cup \bigcup_{\gamma < \beta} C_\gamma$ .

Proof: True if  $\beta = 0$ . Assume for ordinals  $< \beta$ .

(i)  $x \in S_\beta$ .  $\exists$  minimal  $\alpha$ ,  $x \in \varphi_\beta^{-1} \alpha$ . If  $\alpha = 0$  then  $x \in S_\tau$  for some  $\tau < \beta$ .  $\therefore T_x \subseteq S_\tau \cup \bigcup_{\gamma < \tau} C_\gamma \subseteq S_\beta \cup \bigcup_{\gamma < \beta} C_\gamma$ , by 3. If  $\alpha > 0$  then  $\alpha = \tau + 1$  for some  $\tau$ . Since  $x \in \varphi_\beta^{-1} \tau \subseteq \bigcup \varphi_\beta^{-1} \text{On}_\mu$ ,  $T_x \subseteq \bigcup \varphi_\beta^{-1} \text{On}_\mu$  by 11. Let  $y \in T_x$ . If  $y \notin C_\gamma$  for any  $\gamma < \beta$  then  $y \in S_\beta$ ,

by definition of  $S_\alpha : Tx \in S_\alpha \cup C_\gamma$ .

(ii)  $x \in C_\alpha$ . Since  $x \in \varphi_\alpha''On_\alpha$ ,  $Tx \in \varphi_\alpha''On_\alpha$ , by 11. Let  $y \in Tx$ . If  $y \in C_\gamma$  for any  $\gamma < \beta$  then  $y \in S_\beta$  by definition of  $S_\beta : Tx \in S_\beta \cup C_\gamma$ .

13.  $x \in S_\alpha \cup C_\gamma \rightarrow Tx \in S_\alpha \cup C_\gamma$ .

Proof: sufficient to show  $Ux \in S_\alpha \cup C_\gamma$ , and then iterate this result.

Let  $y \in Ux$ .  $\exists z \in x, y \in z$ . By hypothesis,  $z \in S_\alpha \cup C_\gamma$ . If  $z \in S_\alpha$  then  $z \in S_\alpha \cup C_\gamma$ , by 12.  $\therefore y \in S_\alpha \cup C_\gamma$ . Assume  $z \in C_\gamma$ . Then  $\exists$  minimal  $\nu < \beta$ ,  $z \in C_\nu$ .  $\therefore z \in \varphi_\nu''On_\nu$ .  $\therefore y \in \varphi_\nu''\alpha$  for some  $\alpha$ . By definition of  $S_\nu$ ,  $y \in S_\nu$  or  $y \in C_\alpha$  for some  $\alpha < \nu$ .  $\therefore y \in S_\nu \cup C_\nu \in S_\alpha \cup C_\gamma$ .

14.  $x \in \bigcup_{\alpha \in On} S_\alpha \cup Tx \in S_\alpha \cup C_\gamma \rightarrow x \in S_\alpha$ .

Proof: Let  $y \in x$ . If  $x \in S_\alpha$  then  $y \in \varphi_\alpha''On_\alpha$ . If  $y \in C_\gamma$  for some  $\gamma < \alpha$  then  $y \in \varphi_\alpha''0$ . For each  $y \in x$ ,  $\exists$  minimal  $\tau$ ,  $y \in \varphi_\alpha''\tau$ . Since  $x \in \bigcup_{\alpha \in On} S_\alpha$ ,  $x \in S_\alpha$ , so that there is a limit ordinal  $\eta \in On_\alpha$  which is greater than all such  $\tau$ . Then  $y \in x \rightarrow y \in \varphi_\alpha''\tau \in \varphi_\alpha''\eta$ .  $\therefore x \in \varphi_\alpha''\eta$ .  $\therefore x \in \varphi_\alpha''\eta + 1$ .  $\therefore x \in S_\alpha$ .

15.  $x \in \bigcup_{\alpha \in On} S_\alpha \cup Tx \in \bigcup_{\alpha \in On} S_\alpha \cup C_\gamma \rightarrow x \in S_\alpha$ .

Proof: sufficient to show  $\alpha$  minimal,  $x \in S_\alpha \rightarrow \alpha$  minimal,  $Tx \in \bigcup_{\gamma \in On} S_\gamma \cup C_\gamma$ .

(For suppose  $x \in \bigcup_{\gamma \in On} S_\gamma \cup Tx \in \bigcup_{\gamma \in On} S_\gamma \cup C_\gamma$ .  $\exists$  minimal  $\alpha$ ,  $x \in S_\alpha$ . By 12.,  $Tx \in S_\alpha \cup C_\gamma \in \bigcup_{\gamma \in On} S_\gamma \cup C_\gamma$ .  $\therefore Tx \in \bigcup_{\gamma \in On} S_\gamma \cup C_\gamma \rightarrow \alpha \leq \beta$ .  $\therefore x \in S_\alpha$ , by 1.) Use double induction on  $\alpha$  and  $n$ , where  $n$  is defined as follows.  $\exists$  minimal  $\rho > x \in \varphi_\alpha''\rho$ . By minimality of  $\alpha$ ,  $\rho > 0$ . By definition of  $\varphi_\alpha''\rho$ ,  $\sim \text{Lim}(\rho)$ .  $\therefore \exists n, \rho = n + 1$ .  $\therefore x \in \varphi_\alpha''n$ , and  $n$  is minimal with this property.

(i)  $\alpha = 0$ . Then  $Tx \in S_0$ , so  $Tx \in \bigcup_{\gamma \in On} S_\gamma$ .

(ii)  $\alpha = \beta + 1$ .

(a)  $n = 0$ . Then  $x \in S_\alpha \cup C_\beta$ . If  $x \in S_\alpha \cup C_\gamma$ , then  $Tx \in S_\alpha \cup C_\gamma$  by 13.  $\therefore x \in S_\alpha$  by 14.,  $\leftarrow$  minimality of  $\alpha$ .  $\therefore \exists y \in x, y \in C_\beta \wedge y \notin C_\gamma$  for any  $\gamma < \beta$ .  $\therefore Tx \in \bigcup_{\gamma \in On} S_\gamma \cup C_\gamma \rightarrow \beta > \alpha$ , i.e.,  $\beta \geq \alpha$ .

(b)  $n > 0$ . Then  $x \notin \varphi_\alpha''0 = S_\alpha \cup C_\beta$ . But by 12.,  $x \in Tx \in S_\alpha \cup C_\gamma \in S_\alpha \cup C_\beta$ .  $\therefore \exists y \in x, y \in S_\alpha \wedge y \notin S_\gamma$  for any  $\gamma < \beta$ . Now  $y \in \varphi_\alpha''n$ . Since  $\alpha$  is minimal,  $y \in S_\alpha$ ,  $\exists$  minimal  $\delta < \alpha$ ,  $y \in \varphi_\delta''\xi$ . By induction hypothesis,  $\delta < \alpha \rightarrow T_\delta \notin \bigcup_{\gamma \in On} S_\gamma \cup C_\gamma$ .

Since  $\exists y \in Tx, Tx \notin \bigcup_{\gamma < \alpha} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma$ .

(iii)  $\text{Lim}(\alpha)$ .

(a)  $\alpha = 0$ . Then  $x \in \bigcup_{\gamma < \alpha} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma$ . By minimality of  $\alpha, \beta < \alpha \rightarrow x \notin S_\beta \cup C_\beta = \varphi_{\beta+1}^{-1} 0$  (for otherwise  $x \in M \rightarrow x \in \varphi_{\beta+1}^{-1} 1$ , so  $x \in S_{\beta+1}$ ).  $\therefore \exists y \in x$  and  $\delta > \beta, y \in S_\delta \cup C_\delta$  and  $\delta$  is minimal with  $y \in S_\delta \cup C_\delta$ . By induction hypothesis,  $y \in S_\delta \rightarrow \delta$  is minimal,  $Ty \in \bigcup_{\gamma < \delta} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma$ . Since  $Ty \in Tx, Tx \notin \bigcup_{\gamma < \delta} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma$ . If  $y \in C_\delta$ , then  $x \notin \bigcup_{\gamma < \delta} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma$ .  $\therefore Tx \notin \bigcup_{\gamma < \delta} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma$ .

(b)  $\alpha > 0$ . Then  $x \notin \bigcup_{\gamma < \alpha} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma, \beta < \alpha \rightarrow x \notin S_\beta \cup C_\beta$ . Proceed as in case (a).

16.  $x \in \bigcup_{\gamma < \alpha} C_\gamma \wedge Tx \notin \bigcup_{\gamma < \alpha} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma \rightarrow x \in C_\alpha$ .

Proof: Sufficient to show  $\alpha$  minimal,  $x \in C_\alpha \rightarrow \alpha$  minimal,  $Tx \in \bigcup_{\gamma < \alpha} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma$ .

(For suppose  $x \in \bigcup_{\gamma < \alpha} C_\gamma \wedge Tx \notin \bigcup_{\gamma < \alpha} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma$ .  $\exists$  minimal  $\alpha, x \in C_\alpha$ . By 15.,  $Tx \in S_\alpha \cup \bigcup_{\gamma < \alpha} C_\gamma$   $\subseteq \bigcup_{\gamma < \alpha} S_\gamma \cup \bigcup_{\gamma < \alpha} C_\gamma$ .  $\therefore Tx \in \bigcup_{\gamma < \alpha} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma \rightarrow \alpha \leq \beta$ .  $\therefore x \in C_\beta$ , by 5.)

(i)  $\alpha = 0$ . Then  $Tx \in S_0$ , so  $Tx \in \bigcup_{\gamma < \alpha} S_\gamma$ .

(ii)  $\alpha = \beta + 1$ . Required to show  $\exists w \in Tx, w \in C_\beta$  and  $w \notin C_\gamma$  for any  $\gamma < \beta$ . Since  $x \notin \bigcup_{\gamma < \alpha} \varphi_\beta^{-1} 0_{On_\alpha}, \exists y \in x, y \notin \bigcup_{\gamma < \alpha} \varphi_\beta^{-1} 0_{On_\alpha} \therefore y \notin S_\beta$ . By 15.,  $x \in C_{\beta+1} \rightarrow y \in Tx \in S_{\beta+1} \cup \bigcup_{\gamma < \beta} C_\gamma$ .

Suppose  $y \in S_{\beta+1}$ . Then  $\beta + 1$  is minimal with this property, so as in the first sentence of the proof of 15.,  $\exists z \in Ty, z \in C_\beta$  and  $z \notin C_\gamma$  for any  $\gamma < \beta$ .

Since  $z \in Tx, z$  is as required. If  $y \notin S_{\beta+1}$ , then  $y \in C_\beta$ . If  $\exists \gamma < \beta, y \in C_\gamma$ , then  $y \in \varphi_\beta^{-1} 0$ .  $\therefore y \in \bigcup_{\gamma < \alpha} \varphi_\beta^{-1} 0_{On_\alpha}$ ,  $\rightarrow$ . Hence  $y$  is as required.

(iii)  $\text{Lim}(\alpha)$ . For each  $\beta < \alpha$ , required to show  $\exists w \in Tx, w \in C_\beta$  for some minimal  $\delta \geq \beta$ . Since  $x \notin \bigcup_{\gamma < \alpha} \varphi_\beta^{-1} 0_{On_\alpha}, \exists y \in x, y \notin \bigcup_{\gamma < \alpha} \varphi_\beta^{-1} 0_{On_\alpha} \therefore y \notin S_\beta$ . If  $y \in \bigcup_{\gamma < \alpha} S_\gamma$  then  $y \in S_\tau$  for  $\alpha$  minimal  $\tau > \beta$ . As in the first sentence of the proof of 15.,  $Ty \notin \bigcup_{\gamma < \tau} S_\gamma \cup \bigcup_{\gamma < \beta} C_\gamma$ .  $\therefore \exists z \in Ty, z \in C_\delta$  for some  $\delta \geq \beta$  and  $z \notin C_\gamma$  for any  $\gamma < \delta$ . Since  $Ty \in Tx, z$  is as required. Suppose  $y \in \bigcup_{\gamma < \alpha} C_\gamma$ .  $\exists$  minimal  $\nu, y \in C_\nu$ . If  $\nu < \beta$ , then  $y \in \varphi_\beta^{-1} 0$ .  $\therefore y \in \bigcup_{\gamma < \alpha} \varphi_\beta^{-1} 0_{On_\alpha}$ ,  $\rightarrow$ .  $\therefore \nu \geq \beta$ , so  $y$  is as required.

17.  $\alpha \in On \rightarrow M_\alpha = S_\alpha \wedge K_\alpha = C_\alpha$ .

Proof:  $M_\alpha \stackrel{\text{DEFINITION}}{=} \{x \in M : Tx \in M\} \stackrel{9.}{=} \{x \in \bigcup_{\gamma < \alpha} S_\gamma : Tx \in \bigcup_{\gamma < \alpha} S_\gamma\} \stackrel{12, 15}{=} S_\alpha$ .  
 $K_\alpha \stackrel{\text{DEFINITION}}{=} \{x \in V - M : Tx \in M\} \stackrel{9.}{=} \{x \in \bigcup_{\gamma < \alpha} C_\gamma : Tx \in \bigcup_{\gamma < \alpha} S_\gamma\} \stackrel{12, 16}{=} C_\alpha$ . Assume for all  $\nu < \alpha$ .  
 $M_\alpha \stackrel{\text{DEFINITION}}{=} \{x \in M : Tx \in M \cup \bigcup_{\gamma < \alpha} K_\gamma\} \stackrel{9. \text{ INDUCTION HYPOTHESIS}}{=} \dots$

$$\{x \in \bigcup_{\gamma \in \text{On}_M} S_\gamma : \exists \alpha \in \text{On}_M \exists \beta \in \text{On}_M \exists \gamma \in \text{On}_M (x \in S_\gamma \wedge \gamma < \alpha \wedge \gamma < \beta)\} \stackrel{12.15}{=} S_\alpha.$$

$$K_\alpha \stackrel{\text{DEFINITION}}{=} \{x \in V - M : \exists \gamma \in \text{On}_M \exists \delta \in \text{On}_M (x \in V - M \wedge \exists \gamma \in \text{On}_M \exists \delta \in \text{On}_M (x \in V - M \wedge \exists \gamma \in \text{On}_M \exists \delta \in \text{On}_M (x \in V - M \wedge \dots)))\} \stackrel{9.3 \text{ INDUCTION HYPOTHESIS}}{=} \dots$$

$$\{x \in \bigcup_{\gamma \in \text{On}_M} C_\gamma : \exists \alpha \in \text{On}_M \exists \beta \in \text{On}_M \exists \gamma \in \text{On}_M (x \in C_\gamma \wedge \gamma < \alpha \wedge \gamma < \beta)\} \stackrel{12.16}{=} C_\alpha.$$

6. Further properties of K; some cardinality relations.

1. (The following are true also for the system in [1].)

- (a)  $\{x\} \in M \iff \{x_1, \dots, x_n\} \in M$ , by induction.
- (b)  $x \in K_\alpha \implies \{x\} \in M_{\alpha+1}$ .
- (c)  $M \not\subseteq V$ .

Proof: Suppose  $M \subseteq V$ . Since  $M \not\subseteq K$ ,  $M \in C_\alpha$  for some minimal  $\alpha \implies M \in \bigcup_{\gamma < \alpha} \mathcal{P}_\gamma \text{On}_M$ .

Since  $\{C_\alpha\} \in M$ ,  $\{C_\alpha\} \in \mathcal{P}_\beta A$  for some minimal  $\beta$ . Suppose  $\beta = 0$ . Then  $\{C_\alpha\} \in \bigcup_{\gamma < \alpha} S_\gamma \cup C_\gamma \implies \{C_\alpha\} \in S_\gamma$  for some minimal  $\gamma < \alpha \implies \{C_\alpha\} \in \mathcal{P}_\gamma \text{On}_M$  for some minimal  $\gamma$ .  $\implies \{C_\alpha\} \in \mathcal{P}_\gamma \text{On}_M$ , so  $C_\alpha \in \mathcal{P}_\gamma \text{On}_M$ . If  $\gamma = 0$ ,  $C_\alpha \in \bigcup_{\tau < \gamma} S_\tau \cup C_\tau$ . Since  $C_\alpha \not\subseteq S_0$ ,  $C_\alpha \not\subseteq S_\tau \implies C_\alpha \in C_\tau$  for some  $\tau < \alpha$ . But  $\tau < \alpha \implies C_\tau \in C_\alpha \implies \dots \implies \tau \neq 0 \implies \tau = \tau + 1$  for some  $\tau$ . Since  $C_\alpha \in \mathcal{P}_\gamma \text{On}_M$ ,  $C_\alpha \in \mathcal{P}_\gamma \tau \wedge C_\alpha \not\subseteq S_0$ . But  $C_\alpha \not\subseteq S_0 \implies \dots \implies \tau \neq 0 \implies \beta = \gamma + 1$  for some  $\gamma \implies \{C_\alpha\} \in \mathcal{P}_\alpha \text{On}_M \implies C_\alpha \in \mathcal{P}_\alpha \text{On}_M$ . If  $\gamma = 0$ ,  $C_\alpha \in \bigcup_{\tau < \alpha} S_\tau \cup C_\tau$ . If  $C_\alpha \in S_\tau$ ,  $C_\alpha \not\subseteq S_0 \implies \dots \implies C_\alpha \in C_\tau$ . But  $\tau < \alpha \implies C_\tau \in C_\alpha \implies \dots \implies \tau \neq 0 \implies \tau = \tau + 1$  for some  $\tau$ . Since  $C_\alpha \in \mathcal{P}_\alpha \text{On}_M$ ,  $C_\alpha \in \mathcal{P}_\alpha \tau \wedge C_\alpha \not\subseteq S_0 \implies \dots \implies M \not\subseteq V$ .

(d)  $M \cong V$ .

Proof:  $x \mapsto \{x\}$  is 1-1 from  $V$  into  $M$ .

(e)  $x \in \text{On}_M \implies x \in \omega$ .

Proof:  $\text{On}_M \subseteq K_0 \subseteq S_0$ , so  $x \in S_0$ .

2. In the axiom of replacement, if  $x \in M$  then  $y \in M$ .

Proof: Let  $x_1 = x \cdot D(X)$ . Since  $u \in y \iff \exists v \in X, u \in X$ , the correspondence  $v \mapsto u$  maps  $x_1$  onto  $y$ .  $\implies y \in x_1 \in x \in S_0$ .

Notation:  $\alpha_0, \alpha_1, \alpha_2, \dots$  denotes the sequence of ordinal numbers not in  $M$ .

3.  $\beta \in \text{On}_M \implies \alpha_\beta \in C_\beta \wedge (\gamma < \beta \implies \alpha_\gamma \notin C_\gamma)$ .

Proof: (i)  $\beta = 0$ . If  $\gamma < \alpha_0$ , then  $\gamma \in S_0$ , so  $\alpha_\gamma \in S_0$ . Since  $\alpha_0 \notin S_0$ ,  $\alpha_0 \in C_0$ .

Assume for all ordinals  $\alpha < \beta$ .

(ii)  $\beta = \gamma + 1$ . Then  $\alpha_\beta = \alpha_\gamma + 1$ . By induction hypothesis,  $\alpha_\gamma \in C_\gamma$ , so  $\alpha_\gamma \in U \varphi_\gamma "On_M \subseteq U \varphi_{\gamma+1} "On_M$ . Also,  $\alpha_\gamma \in C_\gamma \subseteq \varphi_{\gamma+1}'0 \in U \varphi_{\gamma+1} "On_M$ , so  $\alpha_\beta = \alpha_\gamma \cup \{\alpha_\gamma\} \in U \varphi_{\gamma+1} "On_M \therefore \alpha_\beta \in C_\beta$ . Let  $\delta < \beta$ . To show  $\alpha_\beta \notin C_\delta$ , it is sufficient to show  $\alpha_\beta \notin C_\gamma$ . If  $\alpha_\beta \in C_\gamma$ , then  $\alpha_\gamma \cup \{\alpha_\gamma\} \in U \varphi_\gamma "On_M \therefore \alpha_\gamma \in U \varphi_\gamma "On_M$ . Since by induction hypothesis,  $\alpha_\gamma \notin C_\tau$  for  $\tau < \gamma$ ,  $\alpha_\gamma \in S_\gamma$ ,  $\dashv$ .

(iii)  $\text{Lim}(\beta)$ . Then  $\alpha_\beta = \bigcup_{\gamma < \beta} \alpha_\gamma$ . By induction hypothesis,  $\gamma < \beta \rightarrow \alpha_\gamma \in U \varphi_\gamma "On_M \subseteq U \varphi_\beta "On_M \therefore \alpha_\beta \subseteq U \varphi_\beta "On_M$ . Since  $\gamma < \beta \rightarrow \alpha_\gamma \in S_\gamma$ ,  $\alpha_\beta \notin U S_\gamma \therefore \alpha_\beta \in C_\beta$ . Suppose  $\delta < \beta \wedge \alpha_\beta \in C_\delta$ . Then  $\alpha_{\delta+1} \in \alpha_\beta \subseteq U \varphi_\delta "On_M \therefore \alpha_{\delta+1} \in C_\delta$ ,  $\dashv$ .

4.  $K_\beta \subseteq M_{\beta+1}$ .

Proof:  $x \mapsto \{x\}$  is 1-1 from  $C_\beta$  into  $S_{\beta+1}$ .

$\supset$ .  $x \in V-M \rightarrow \{u : u \in x \wedge u \in S_\alpha\} = x$ .

Proof: Since  $x \subseteq S_\alpha$ ,  $x \cap V-M \subseteq S_\alpha \rightarrow x \approx x$ , by induction on  $\alpha$ .  $\therefore \{u : u \in x \wedge u \in S_\alpha\} = x$ .

$\subseteq U \{u : u \in x \wedge u = \alpha\} \cup \{x\} \subseteq S_\alpha = x$ .

6.  $S_\alpha \subseteq U \varphi_\alpha "On_M$ .

Proof:  $S_\alpha = U \varphi_\alpha "On_M$ . Let  $\alpha > 0$ . For all  $\beta$ ,  $C_\beta \cap S_\alpha$ , so if  $\alpha > 0$ ,  $\varphi_\alpha'0 = \bigcup_{\gamma < \alpha} S_\gamma \cup C_\gamma \cap S_\alpha \therefore \varphi_\alpha'0 \in V-M$ . Assume  $\varphi_\alpha' \delta \in V-M$  for all  $\delta < \alpha$ . If  $\alpha = \beta + 1$ ,  $\varphi_\alpha' \alpha = \varphi_\beta' \beta$  by  $\supset$ .

$\therefore \varphi_\alpha' \alpha \in V-M$ . If  $\text{Lim}(\alpha)$ ,  $\varphi_\alpha' \alpha = U \varphi_\alpha " \alpha$ , so  $\varphi_\alpha' \alpha \in V-M$ . Now  $\varphi_\alpha'0 = \bigcup_{\gamma < \alpha} S_\gamma \cup C_\gamma \approx U C_\gamma$ , since  $C_\gamma \supset S_\gamma$ . Assume  $\varphi_\alpha' \delta \notin U C_\gamma$  for all  $\delta < \alpha$ . If  $\alpha = \beta + 1$ ,  $\varphi_\alpha' \alpha = \varphi_\beta' \beta$  by  $\supset$ , so  $\varphi_\alpha' \alpha \in U C_\gamma$ . If  $\text{Lim}(\alpha)$ ,  $\varphi_\alpha' \alpha = \bigcup_{\gamma < \alpha} \varphi_\alpha " \gamma = \bigcup_{\gamma < \alpha} (U C_\gamma) \cap \alpha = \bigcup_{\gamma < \alpha} C_\gamma$ .  $\therefore \bigcup_{\gamma < \alpha} \varphi_\alpha " On_M \approx (U C_\gamma) \cap On_M \approx (U C_\gamma) \cap (On_M - \{0\}) \approx U \varphi_\alpha "(On_M - \{0\}) \approx \bigcup_{\gamma < \alpha} \varphi_\alpha " On_M \cup C_\gamma = S_\alpha$ .

7.  $x \in C_\alpha \rightarrow x \subseteq S_\alpha$ .

Proof:  $x \subseteq U \varphi_\alpha " On_M \subseteq S_\alpha$  by 6.

8.  $C_\alpha = P(S_\alpha) - S_\alpha \approx PS_\alpha$ .

Proof:  $C_\alpha = P(U \varphi_\alpha " On_M) - S_\alpha \approx PS_\alpha - S_\alpha$  by 6. But  $PS_\alpha \supset S_\alpha$ .

9.  $C_\alpha \approx S_{\alpha+1}$ .

Proof:  $\varphi_{\alpha+1}'0 \in C_\alpha$ . For all  $\beta \in On_M$ ,  $\varphi_{\alpha+1}'\beta \in C_\alpha$  by  $\supset$ , and induction on  $\beta$ .

$\therefore U \varphi_{\alpha+1}' On_M \subseteq C_\alpha \times S_\alpha \subseteq C_\alpha$ . But  $S_{\alpha+1} \approx U \varphi_{\alpha+1}' On_M$  by 6.

10.  $\alpha < \beta \rightarrow S_\alpha \subseteq S_\beta$ .

Proof:  $\alpha + 1 \leq \beta$ , so  $S_\alpha \subseteq S_{\alpha+1} \subseteq S_\beta$ . But  $S_\alpha \subseteq S_{\alpha+1}$  by 9.

11.  $\forall \alpha < \kappa, \exists S_\alpha$ .

Proof:  $\forall \alpha < \kappa$  by induction on  $\alpha$ . By 3.,  $\alpha_0 \in C_\kappa$ , so by 7.,  $\alpha_0 < S_\kappa$ .

12.  $\gamma$  is not an initial ordinal  $\rightarrow \forall \alpha < S_\gamma$ .

Proof:  $\exists$  initial ordinal  $\beta < \gamma, \beta \approx \gamma$ . By 11.,  $\beta \in S_\beta$ .  $\therefore \gamma \approx \beta \in S_\beta \wedge S_\gamma$ , by 10.

13.  $\text{Lim}(\alpha) \rightarrow S_\alpha \approx \bigcup_{\beta < \alpha} S_\beta$ .

Proof:  $\forall S_\alpha \approx \bigcup_{\beta < \alpha} C_\beta$  since  $S_\beta \wedge C_\beta \approx S_{\beta+1}$  for all  $\beta$ . By the proof of 6.,

$\forall \alpha' \beta \approx \bigcup_{\gamma < \alpha'} C_\gamma$  for all  $\beta \in \text{On}_M$ .  $\therefore \forall \alpha' \beta \in \text{On}_M \approx (\bigcup_{\gamma < \alpha'} C_\gamma) \wedge S_\beta \approx \bigcup_{\gamma < \alpha'} C_\gamma$ .  $\therefore S_\alpha \approx \bigcup_{\beta < \alpha} C_\beta$  by 6.

14.  $\gamma$  not an initial ordinal and  $\text{Lim}(\gamma) \rightarrow \bar{S}_\gamma$  not regular.

Proof:  $\exists$  a smaller (than  $\bar{S}_\gamma$ ) increasing sequence of smaller (than  $\bar{S}_\gamma$ ) cardinal numbers whose union is  $\geq \bar{S}_\gamma$ , viz.,  $\{\bar{S}_\alpha : \alpha < \gamma\}$ , by 10., 12. and 13.

15. If  $\text{Lim}(\gamma)$ ,  $\bar{S}_\gamma \approx \bigcup_{\alpha < \gamma} \bar{S}_\alpha$ . If  $\bar{S}_\gamma$  is also regular,  $S_\gamma \approx \gamma$ .

Proof: By 13.,  $\bar{S}_\gamma \approx \bigcup_{\alpha < \gamma} \bar{S}_\alpha$ , and by 10.,  $\bar{S}_\alpha \wedge \bar{S}_\gamma \approx \bar{S}_\alpha \wedge S_\gamma \approx S_\gamma$  by 11. Hence, if  $\bar{S}_\gamma$  is regular,  $\bar{S}_\gamma \not\approx \bar{S}_\gamma$ .  $\therefore S_\gamma \approx \gamma$ , by 11.

16.  $x \wedge S_\alpha \wedge \text{Lim}(\alpha) \rightarrow \text{Px} \wedge S_\alpha$ .

Proof: If  $x \approx S_\gamma$  for all  $\gamma < \alpha$  then  $x \approx \bigcup_{\gamma < \alpha} S_\gamma$ .  $\therefore x \approx S_\alpha$  by 13.,  $\rightarrow$ .  $\therefore x \approx S_\gamma$  for some  $\gamma < \alpha$ .  $\therefore \bar{S}_\gamma \approx \bar{S}_\gamma \approx \bar{S}_\gamma \wedge \bar{S}_\alpha \approx \bar{S}_\alpha$  by 6., 9., and 10.

17.  $\forall \alpha \in \text{On}_M \rightarrow \bar{S}_\alpha$  not regular or not  $\text{Lim}(\alpha)$ .

Proof: Otherwise  $\bar{S}_\alpha \approx \bar{S}_\alpha$  by 15., but  $\bar{S}_\alpha \notin \text{On}_M$ .

18.  $\forall \alpha \in \text{On}_M \rightarrow \alpha_\gamma \approx S_0$ .

Proof: By 3. and 7.,  $\alpha_0 \approx S_0$ . Assume for ordinals  $\alpha < \gamma$ . If  $\gamma \approx \beta+1$ ,  $\alpha_\gamma \approx \alpha_\beta + 1 \approx \alpha_\beta \approx S_0$ . If  $\text{Lim}(\gamma)$ ,  $\alpha_\gamma \approx \bigcup_{\beta < \gamma} \alpha_\beta$ , so  $\bar{\alpha}_\gamma \approx \bar{S}_0 \approx \bar{S}_0$ .

19.  $\forall \alpha \in \text{On}_M \rightarrow \bar{S}_{\alpha_\gamma}$  not regular or not  $\text{Lim}(\alpha_\gamma)$ .

Proof: Otherwise  $\bar{S}_{\alpha_\gamma} \approx \bar{S}_{\alpha_\gamma}$  by 15. But  $\bar{\alpha}_\gamma \approx \bar{S}_0$  by 18.  $\therefore \bar{S}_{\alpha_\gamma} \approx \bar{S}_0$ ,  $\rightarrow$  10.

20.  $\bar{S}_0$  is regular. (By 18.,  $\bar{S}_0 \approx \alpha_0$ . Also  $\alpha_0 \approx \omega_{\alpha_0}$  and  $\text{Lim}(\alpha_0)$ . Hence  $\alpha_0$  is inaccessible;  $\therefore \text{PS}_0$  is a model for NBG by [5].)

Proof: Let  $z$  be an increasing sequence of cardinal numbers,  $D(z) \wedge S_0 \wedge (y \in R(z) \rightarrow y \wedge S_0)$ . Then  $R(z) \wedge M \wedge R(z) \approx M$ .  $\therefore \bigcup R(z) \approx M \wedge R(z) \wedge M$ .  $\therefore \bigcup R(z) \wedge S_0$ .

Notation:  $\alpha_0, \alpha_1, \alpha_2, \dots$  denotes the sequence of initial ordinals not in  $M$ .

21. If the generalized continuum hypothesis (GCH) holds for the  $\aleph_\alpha$ , then  $\overline{S}_\alpha = \aleph_\alpha$  for all  $\alpha$ .

Proof:  $\overline{S}_0 = \aleph_0$ . Assume for all ordinals  $\alpha < \beta$ . Let  $\beta = \gamma + 1$ . Then  $\overline{S}_\gamma = \aleph_\gamma$ , so  $\aleph_{\gamma+1} = \aleph_\gamma^{\aleph_\gamma} = 2^{\aleph_\gamma} = \overline{S}_{\gamma+1}$ , by 8. and 9. Assume  $\text{Lim}(\beta)$ . By 15.,  $\overline{S}_\beta = \bigcup_{\gamma < \beta} \overline{S}_\gamma = \bigcup_{\gamma < \beta} \aleph_\gamma = \aleph_\beta$ .

22. If  $\overline{S}_\alpha = \aleph_\alpha$  for all  $\alpha$ , then GCH holds for the  $\aleph_\alpha$ .

Proof:  $\aleph_S \approx S_{\alpha+1}$  by 8. and 9.

23.  $\tau$  minimal,  $\forall \delta \leq S_\tau \rightarrow \tau$  minimal,  $\aleph_\tau \leq S_\tau$ . (Hence if  $\gamma \neq S_\tau$  then  $S_\gamma \not\leq \aleph_\gamma$  by 11., so  $\aleph_\gamma \geq S_\tau$ .)

Proof: Show first  $\aleph_\tau \leq S_\tau$ . True if  $\tau = 0$ , by 10. Assume  $\tau = \beta + 1$ . Then  $\tau \neq \beta$ , so  $\tau \neq S_\tau \rightarrow \beta \leq S_\tau$ . Also,  $\tau$  is minimal,  $\beta \leq S_\tau$ ; hence  $\aleph_\beta \leq S_\tau$ . But  $\aleph_{\beta+1} = \aleph_\beta + 1 \not\leq \aleph_\beta$ . Assume  $\text{Lim}(\tau)$ . Then  $\aleph_\tau = \bigcup_{\delta < \tau} \aleph_\delta$ . Since  $\forall \delta \leq S_\tau, \delta < \tau \rightarrow \delta \leq S_{\tau(\delta)}$  for some minimal  $\tau(\delta) \leq \delta$ .  $\therefore \aleph_\tau \leq \bigcup_{\delta < \tau} S_{\tau(\delta)} \leq S_\tau$ . If  $\beta < \tau$  and  $\aleph_\beta \leq S_\tau$  then  $\tau \leq S_\beta$  by 11.,  $\rightarrow$ .

24.  $\beta \neq \gamma \rightarrow \aleph_\beta \not\leq \aleph_\gamma$ .

Proof: Use induction on  $\gamma$ . Assume  $\beta < \gamma$ . True if  $\gamma = 0$ . Suppose  $\gamma = \delta + 1$ . Then  $\beta \leq \delta \wedge \beta \neq \delta \rightarrow \aleph_\beta \not\leq \aleph_\delta$ . But  $\aleph_\gamma = \aleph_\delta + 1 \not\leq \aleph_\delta$ . Suppose  $\text{Lim}(\gamma)$ . Then  $\aleph_\gamma = \bigcup_{\tau < \gamma} \aleph_\tau = \bigcup_{\beta < \tau < \gamma} \aleph_\tau$ . But  $\beta < \tau < \gamma \wedge \beta \neq \tau \rightarrow \beta \neq \tau$ .  $\therefore \aleph_\tau \not\leq \aleph_\beta$ .  $\therefore \aleph_\gamma \not\leq \aleph_\beta$ .  $\overline{S}_\gamma = \overline{S}_\beta \cdot \overline{S}_\gamma = \overline{S}_\beta$ .

25.  $(\exists \omega_n \text{ is On}_n) \left( \gamma = \sum_{i=0}^n \omega_i \right) \rightarrow \gamma \not\leq \omega_0$ .

Proof: Use induction on  $n$ . True for  $n=0$ , by 15. Let  $\beta = \sum_{i=0}^n \omega_i$ . By

induction hypothesis,  $\beta \not\leq \omega_0$ .  $\therefore \aleph_\beta \not\leq \aleph_{\omega_0}$  by 24. But  $\aleph_{\omega_0} = \bigcup_{\tau < \omega_0} \aleph_\tau$ , so  $\overline{S}_{\omega_0} \leq \overline{S}_\beta \cdot \overline{S}_{\omega_0} = \overline{S}_\beta$ .

26.  $\aleph_{\aleph_\gamma} = \aleph_\gamma$  for all  $\gamma$ .

Proof: Suppose  $\gamma$  is the least ordinal,  $\aleph_{\aleph_\gamma} \neq \aleph_\gamma$ . Then  $\aleph_{\aleph_\gamma} > \aleph_\gamma$  by 11. Now  $\aleph_\gamma = \aleph_\delta$  for some  $\delta$ , and  $\delta \leq \aleph_\beta$  for some  $\beta$ .  $\therefore \aleph_\gamma = \aleph_\delta \not\leq \aleph_{\aleph_\beta}$  by 24. Since  $\aleph_{\aleph_\gamma} > \aleph_{\aleph_\beta}$ ,  $\aleph_\gamma > \aleph_\beta$ . Hence  $\gamma > \beta$ .  $\therefore \aleph_\gamma \not\leq \aleph_\beta$ .  $\therefore \aleph_\gamma \not\leq \aleph_\beta$ .  $\therefore \gamma = \beta$ ,  $\rightarrow$ .

27.  $\aleph_\gamma = \gamma \wedge (\text{EA})(\beta \leq \delta < \gamma \rightarrow \delta < \aleph_\delta) \rightarrow \gamma$  is not regular. (This holds in general for initial ordinals  $\omega_\gamma$ .)

Proof: Let  $\beta \leq \delta < \gamma$ . Then  $\delta < \aleph_\delta < \aleph_\gamma = \gamma$ .  $\therefore \beta < \aleph_\delta < \gamma \rightarrow \aleph_\beta < \aleph_{\aleph_\delta} < \aleph_\gamma = \gamma$ . By iteration,

$\alpha < \gamma$ . Suppose  $\alpha < \gamma$  and  $\alpha = \text{l.u.b.} \{ \alpha_i : i \in \omega \}$ . For some  $\rho, \tau \leq \rho$ , and  $\rho \leq \rho$  by induction. Since  $\alpha_i < \rho, \alpha_i < \rho$  for all  $i \in \omega, \therefore \rho \geq \alpha, \therefore \rho = \alpha$ , i.e.,

$\alpha < \alpha$ . But  $\alpha < \alpha, \alpha < \alpha$ , so  $\alpha < \alpha, \therefore \alpha = \text{l.u.b.} \{ \alpha_i : i \in \omega \}$ , i.e.,  $\alpha$  is the limit of a smaller (than  $\alpha$ ) increasing sequence of smaller (than  $\alpha$ ) ordinals.

7. Supercomplete inner models of K.

Let  $G$  denote the von Neumann function. Following [5], define multiplicity, extended class and extended set, respectively, of a supercomplete inner model of  $K$  by  $X \leq W, x \leq W, x \leq W \cdot M$ , where  $W$  is a multiplicity of  $K, X$  is a variable ranging over the multiplicities of  $K$ , and  $x$  is a variable ranging over the extended classes and extended sets of  $K$ . Since any model for  $K$  is a model for NBG, either  $W = V$  or  $W = G' \cdot \gamma$  where  $\gamma$  is an inaccessible ordinal of  $K$ . If  $\gamma \leq M$ , then the model becomes an interpretation of  $K$  with the extended classes the same as the extended sets, since  $G' \cdot \gamma \leq M$ . Hence, assume  $\gamma = \alpha_\rho$  for some  $\rho$ .

1. If  $z \in G' \cdot \alpha_\rho$  and  $\gamma$  is minimal,  $z \in M_\gamma \cup K_\gamma$ , then  $\gamma \leq \beta$ . (Suppose  $\text{Lim}(\gamma)$  and  $z \in G' \cdot \alpha_\tau$ . Then  $z \in G' \cdot \alpha_\beta$  for some  $\beta < \tau$ . Hence if  $\gamma$  is minimal,  $z \in M_\gamma \cup K_\gamma, \gamma \leq \beta < \tau$ .)

Proof: True for  $\beta = 0$ , since  $G' \cdot \alpha_0 = M_0$ . Assume for ordinals  $< \beta$ . Suppose  $\beta = \delta + 1$ . Then  $z \in G' \cdot \alpha_\delta$ . Suppose  $\exists \eta_1 < \gamma \wedge \eta_2 < \gamma, z \in M_{\eta_1} \cup K_{\eta_1}$ . Then  $\exists z \in M_{\eta_1} \cup K_{\eta_1} \subseteq M_{\eta_2} \cup K_{\eta_2}$  where  $\eta = \max(\eta_1, \eta_2)$ .  $\therefore z \in M_\eta \cup K_\eta$ ,  $\leftarrow$  minimality of  $\gamma$ . Hence either  $\exists y \leq z, \gamma$  is minimal with  $y \in M_\gamma$ , or  $\eta < \gamma \rightarrow \exists y \leq z, y \in K_\eta$  for some minimal  $\rho$  with  $\eta < \rho < \gamma$ . Now  $y \in G' \cdot \alpha_\rho$ . Hence by induction hypothesis, in the first case  $\gamma \leq \delta < \beta$ , and in the second case  $\eta < \rho \leq \delta < \beta$ , so that  $\gamma \leq \beta$ . Suppose  $\text{Lim}(\beta)$ . Then  $z \in G' \cdot \alpha_\tau$  for some  $\tau < \beta, \therefore \gamma \leq \tau < \beta$ .

2. If  $\kappa < \beta$ , then  $M_\kappa \leq G' \cdot \alpha_\beta$ . (Suppose  $\alpha_\tau$  is inaccessible and  $\kappa < \tau$ . Let  $\kappa < \beta < \tau$ . Then  $M_\kappa \leq G' \cdot \alpha_\beta$ . Since  $G' \cdot \alpha_\beta \leq G' \cdot \alpha_{\beta+1} \leq G' \cdot \alpha_\tau, G' \cdot \alpha_\beta \leq G' \cdot \alpha_\tau$ , so  $M_\kappa \leq G' \cdot \alpha_\tau$ .)

Proof: True for  $\beta = 0$ . Assume for ordinals  $< \beta$ . Suppose  $\text{Lim}(\beta)$ . Let  $\kappa < \tau < \beta$ . Then  $M_\kappa \leq G' \cdot \alpha_\tau \leq G' \cdot \alpha_\beta$ . Suppose  $\beta = \tau + 1$ . If  $\kappa < \tau$  then  $M_\kappa \leq G' \cdot \alpha_\tau \leq G' \cdot \alpha_{\tau+1} = G' \cdot \alpha_{\beta}$ . Hence

it suffices to show  $M_\tau \models G'_{\alpha_{\tau+1}} \equiv P(G'_{\alpha_\tau})$ . True for  $\tau=0$ . Assume for ordinals  $\alpha_\tau$ . If  $\tau = \gamma+1$ ,  $M_\gamma \models G'_{\alpha_\tau}$ , so  $M_\tau \models P(M_\gamma) \models P(G'_{\alpha_\tau})$ . If  $\text{Lim}(\tau)$ ,  $\overline{M}_\tau = \bigcup_{\tau \in C} \overline{M}_\tau \models \bigcup_{\tau \in C} G'_{\alpha_{\tau+1}} \equiv \tau \cdot G'_{\alpha_\tau} \equiv G'_{\alpha_\tau}$ , since  $\tau \leq \alpha_\tau$  for all  $\tau$  and  $\beta \leq G'_{\alpha_\tau}$  by induction.

3. If  $\alpha_\rho$  is an inaccessible ordinal of  $K$  and  $\rho > 0$ , then  $G'_{\alpha_\rho}$  determines a supercomplete inner model of  $K$ .

Proof: Since  $G'_{\alpha_\rho}$  is a model for NBG it is necessary only to check the special axioms. In the following proofs of the special axioms, let  $(*)$  denote the fact that the special axiom in question holds for the universe  $V$  of  $K$ , and  $(**)$  the fact that  $G'_{\alpha_\rho}$  is a model for NBG.

(a) Subsets $_{\mathfrak{M}}$ : Let  $x \in M \cdot G'_{\alpha_\rho}$ ,  $Z \subseteq G'_{\alpha_\rho}$ . Then  $x \cdot Z \in M$  by  $(*)$  and  $x \cdot Z \in G'_{\alpha_\rho}$  by  $(**)$ .

(b) Union $_{\mathfrak{M}}$ : Let  $x \in M \cdot G'_{\alpha_\rho}$ . Then  $x \cdot G'_{\alpha_\rho} \in M$  by Subsets $_{\mathfrak{M}}$  for the system  $K$ , so  $\bigcup x \cdot M \cdot G'_{\alpha_\rho} \in M$  by  $(*)$ . Also,  $x \cdot M \cdot G'_{\alpha_\rho} \in G'_{\alpha_\rho}$  by  $(**)$ , so  $\bigcup x \cdot M \cdot G'_{\alpha_\rho} \in G'_{\alpha_\rho}$  by  $(**)$ .

(c) Power $_{\mathfrak{M}}$ : Let  $x \in M \cdot G'_{\alpha_\rho}$ . Then  $Px \in M$  by  $(*)$  and  $Px \in G'_{\alpha_\rho}$  by  $(**)$ .

(d) Pairing $_{\mathfrak{M}}$ : Let  $x, y \in M \cdot G'_{\alpha_\rho}$ . Then  $\{x, y\} \in M$  by  $(*)$  and  $\{x, y\} \in G'_{\alpha_\rho}$  by  $(**)$ .

(e) Infinity $_{\mathfrak{M}}$ : Let  $x$  be given by Infinity $_{\mathfrak{M}}$  for  $K$ . Then  $x \in M_0 = G'_{\alpha_0} \subseteq G'_{\alpha_\rho}$ .  $\therefore Tx \in G'_{\alpha_\rho}$ .

(f) Cardinality: If  $x \in M \cdot G'_{\alpha_\rho}$ ,  $\exists y \in M_0, x \approx y$ , by  $(*)$ .  $\therefore y \in G'_{\alpha_0} \subseteq G'_{\alpha_\rho}$ . The converse, viz.,  $(x \in G'_{\alpha_\rho} \wedge \exists y \in M_0, x \approx y \rightarrow x \in M)$ , also follows from  $(*)$ .

(g) Universes: Let  $z \in G'_{\alpha_\rho}$ . Required to show  $y = \{x \in G'_{\alpha_\rho} : Tx \in M \cdot G'_{\alpha_\rho} \cup z\} \in G'_{\alpha_\rho}$ . Sufficient to show  $y \in G'_{\alpha_\rho}$ , for then  $y \in \alpha_\rho$ , so  $y \neq \beta$  for some  $\beta \in G'_{\alpha_\rho}$ , and the axiom of replacement applies.  $\exists$  minimal  $\alpha, z \in \alpha \cup K_\alpha$ . Then  $z \in M_\alpha \cup K_\alpha$ .  $\therefore y \in \{x : Tx \in M \cup K_\alpha\} = M_{\alpha+1} \cup K_{\alpha+1} \subseteq K_{\alpha+2}$ . Hence to show  $y \in G'_{\alpha_\rho}$  it suffices to show  $M_{\alpha+2} \in G'_{\alpha_\rho}$ . By 1.,  $\alpha < \rho$ , so  $\alpha+2 < \rho$ . By 2.,  $M_{\alpha+2} \in G'_{\alpha_\rho}$ .

### 3. Relative consistency.

An interpretation of  $K$  is constructed in  $\text{NBG} + (\text{AI})$ , where AI is the axiom of inaccessibility. This is done in two ways, with A. much simpler than B. (Apparently B. is superfluous.) Note that for A. the underlying set theory may be taken to be  $\text{NBG} + (\text{E } 2 \text{ inaccessible ordinals})$ ; the

assumption AI is not necessary. Let  $G$  denote the von Neumann function.  
 A. For any  $\delta$  let  $N_\delta = \{x: x < \delta\}$ . If  $\gamma$  is minimal,  $x \in G'\gamma$ , call  $\gamma$  the index of  $x$ .  
 1. If  $\rho$  is a regular limit ordinal and  $Tx \in N_\rho$ , then  $x \in G'\rho$ .

Proof: Use induction on the index of  $x$ . True for index 1; assume for indices smaller than that of  $x$  and let  $x$  have index  $\alpha+1$ . Let  $y \in x$  and let  $\beta+1$  be the index of  $y$ . Since  $Ty \in Tx \in N_\rho$  and  $\beta+1 < \alpha+1$ ,  $y \in G'\rho$ . Also,  $y < \rho$ .  $\therefore$  the indices of the elements of  $y$  can be arranged in an increasing sequence  $(\beta_\nu)_{\nu < \kappa}$  for some  $\kappa < \rho$ , where  $\beta_\nu < \beta$  for all  $\nu < \kappa$ . By regularity,  $\bigcup_{\nu < \kappa} \beta_\nu < \rho$ . But  $\beta$  is minimal,  $y \in G'\beta$ .  $\therefore$  if  $\nu < \beta$ ,  $\exists z \in y, z \notin G'\nu$ . If the index of  $z$  is  $\nu < \nu$ , then  $z \in G'\nu \in G'\nu$ .  $\therefore \nu > \nu$ .  $\therefore \bigcup_{\nu < \kappa} \beta_\nu \geq \beta$ .  $\therefore \beta < \rho$ .  $\therefore y \in G'(\beta+1) \in G'\rho$ .

2. Assume  $\beta$  inaccessible. (Also, assume  $\gamma$  inaccessible with  $\gamma > \beta$ .) Then  $V(G'\gamma)$  determines a model for  $K$ , where multiplicity, extended class and extended set of one model, respectively, are defined in  $NBG+(AI)$  by  $x \in V$  ( $x \in G'\gamma$ ),  $x \in V$  ( $x \in G'\gamma$ ),  $x \in N_A$  ( $x \in G'\gamma \cdot N_A$ ), where  $x$  is a class and  $x$  is a set in the sense of  $NBG+(AI)$ .

Note: These are the only possible interpretations of  $K$  in  $NBG+(AI)$  with the notions of multiplicity, extended class and extended set defined by  $x \in W$ ,  $x \in U$ ,  $x \in V$  for some classes  $W, U$  of  $NBG+(AI)$ ; for in any such interpretation the universe of  $K$  must satisfy the axioms of  $NBG$  (hence must be  $V$  or  $G'\gamma$ ),  $M_0$  must satisfy the axioms for  $NBG$  (hence must be  $G'\beta$ ), and  $V$  must consist of the sets cardinally less than  $M_0$ . Also, if  $K$  has such an interpretation in an extension  $NBG\#$  of  $NBG$ , the  $M_0$  of the interpretation is a set in  $NBG\#$  which determines a model for  $NBG$ , hence must be  $G'\beta$ , where  $\beta$  is an inaccessible ordinal in  $NBG\#$ .

Proof: Since  $V(G'\gamma)$  satisfies the axioms for  $NBG$ , it is necessary only to check the special axioms. All of these are immediate but Universes and perhaps Union<sub>m</sub>. For Union<sub>m</sub>, let  $x \in N_\beta$ . Then  $\overline{\bigcup_{z \in x} z} = \overline{\bigcup_{z \in x} z}$ . Since  $\overline{x \cdot N_\beta} \in \overline{x} < \beta$ ,  $\bigcup_{z \in x} \overline{z}$  is the union of a smaller (than  $\beta$ ) set of smaller (than  $\beta$ ) cardinals, hence is  $< \beta$ . For Universes, let  $z \in V$  ( $z \in G'\gamma$ ) with index  $\delta+1$ . Then  $\{x: Tx \in N_\delta \cdot z\} \in$

$\{x: Tx \in N_{\bar{\alpha}} \cup G'\bar{\alpha}\} \subseteq \{x: Tx \in N_{\tau}\} \quad (\{x \in G'\gamma: Tx \in N_{\bar{\alpha}} \cup Z\} \subseteq \{x \in G'\gamma: Tx \in N_{\bar{\alpha}} \cup G'\bar{\alpha}\} \subseteq \{x \in G'\gamma: Tx \in N_{\tau}\})$   
 where  $\bar{\tau} = \max(\bar{\alpha}, \overline{G'\bar{\alpha}})$ .  $\bar{\tau}$  regular limit ordinal  $\rho, \tau < \rho$  ( $\tau < \rho < \gamma$ ). Then by 1.,  
 $\{x: Tx \in N_{\tau}\} \subseteq \{x: Tx \in N_{\rho}\} \subseteq G'\rho + 1 \quad (\{x \in G'\gamma: Tx \in N_{\tau}\} \subseteq \{x \in G'\gamma: Tx \in N_{\rho}\} \subseteq G'\rho + 1)$ .  
 $\therefore \{x: Tx \in N_{\bar{\alpha}} \cup Z\} \in V \quad (\{x \in G'\gamma: Tx \in N_{\bar{\alpha}} \cup Z\} \in G'\gamma)$ .

B. Let  $\tau$  be an increasing sequence of inaccessible ordinals,

$$R_{\tau_n} = G'\tau_n, \quad M_n = R_{\tau_n}, \quad K_n = Pn_{\tau_n} - M_n,$$

$$M_n = \{x \in R_{\tau_n} : (E y \in M_n, x \neq y) \wedge (\forall z \in Tx \rightarrow (E w \in M_n, y = w) \vee \exists_{\tau_n} U K_n))\},$$

$$K_n = \{x \in Pn_{\tau_n} - M_n : \forall z \in Tx \rightarrow (E w \in M_n, y = w) \vee \exists_{\tau_n} U K_n\}.$$

Use the notation  $W = \bigcup_{n \in \omega} M_n \cup K_n$  for the universe of the interpretation,  $(\Phi_n)$

for the statement that  $R_{\tau_n}$  is a model for NBG,  $(\Phi_n)$  for the statement

$(\forall z \in Tx \rightarrow (E w \in M_n, y = w) \vee \exists_{\tau_n} U K_n)$ ,  $x, y, \dots$  for elements of  $W$ , and  $X, Y, \dots$  for

subclasses of  $W$ .

Note: Since this interpretation of  $K$  is of the type described in A., the dependence of  $M_n$  and  $K_n$  upon the particular sequence  $(\tau_n)$  is only apparent, and the interpretation is determined by  $\tau_n$ . Hence by definition of  $M_n$  and  $K_n$ ,  $x \in M_n \rightarrow x \in R_{\tau_n} \wedge x \in K_n \rightarrow x \in R_{\tau_n}$  for any increasing sequence  $\tau$  of inaccessible ordinals beginning with  $\tau_0$ , in particular, for the sequence of all such.

$$1. \delta \leq \beta \rightarrow M_\delta \subseteq M_\beta.$$

$$\text{Proof: } x \in M_\delta \rightarrow x \in R_{\tau_\delta} \subseteq R_{\tau_\beta}.$$

$$2. K_n \cdot M_\nu = \emptyset.$$

Proof: By definition of  $K_n$ ,  $K_n \cdot M_n = \emptyset$ .  $\therefore \forall \delta \leq \beta \rightarrow K_n \cdot M_\nu \subseteq K_n \cdot M_n = \emptyset$ , by 1. Suppose

$\beta < \nu$ . If  $\nu = 0$ ,  $K_n \cdot M_\nu = \emptyset$ . Assume for ordinals  $\lambda < \nu$ . Let  $\nu = \delta + 1$ . If  $\beta < \delta$  then

$K_n \cdot M_\delta = \emptyset$  by induction hypothesis. Let  $\beta < \delta$  and assume  $x \in K_n \cdot M_{\delta+1}$ . Since

$$x \in K_n, x \in M_n \rightarrow x \in R_{\tau_n} \wedge \exists y \in M_n, x \neq y. \text{ But } x \in M_{\delta+1} \rightarrow x \in R_{\tau_n}. \text{ Since } K_n \cdot M_\delta = \emptyset, x \notin M_\delta.$$

$\therefore x \notin R_{\tau_n} \rightarrow x \in R_{\tau_{\delta+1}} - R_{\tau_n}$ .  $\exists$  minimal  $n$ ,  $x \in G'n + 1$  and  $\tau_\delta \leq n < \tau_{\delta+1}$ . Since  $x \in K_n$ ,

$$x \in G'\tau_n + 1. \text{ But } \tau_n < \tau_\delta \leq n, \rightarrow \text{minimality of } n. \therefore \beta < \delta \rightarrow K_n \cdot M_{\delta+1} = \emptyset. \text{ Suppose}$$

$x \in K_n \cdot M_{\delta+1}$ . Then  $x \in R_{\tau_{\delta+1}} - R_{\tau_\delta}$ , since  $K_n \cdot M_\delta = \emptyset$ .  $\exists y \in M_n, x \neq y$ , and  $\exists$  minimal

$n < \tau_n, y \in G'n + 1$ , i.e.,  $y \in G'n$ . If  $n \neq \delta$ , then  $\exists$  minimal  $\rho \leq n, w \in G'\rho$ . Since

$x \in G'\tau_i$ , the set of corresponding  $\rho$  for all  $w \in x$  forms a sequence with domain  $\subseteq G'\tau_i$ ; but  $G'w \in G'\tau_i$ , so by  $(\Phi_i)$ ,  $G'w \in \tau_i$ . Also, these  $\rho$  are all  $\in \tau_i$ . By regularity of  $\tau_i$ , the union  $\xi$  of these  $\rho$  is also  $\in \tau_i$ . Assume  $\text{Lim}(\mu) \wedge \xi \subseteq \mu \in \tau_i$ . If  $w \in x$ ,  $\exists$  minimal  $\alpha, w \in G'\alpha$ . Since  $G'\alpha \in G'\mu$ ,  $x \in G'\mu$ .  $\therefore x \in G'\mu + 1 \in R_{\tau_i}$ ,  $\therefore K_{\beta} \cdot M_{\beta+1} = \emptyset$ . Assume  $\text{Lim}(\gamma), \beta < \gamma, x \in K_{\beta} \cdot M_{\beta}$ . Since  $x \in K_{\beta}$ ,  $x \in G'\tau_{\beta}$ .  $\therefore x \in G'\tau_{\beta} + 1 \in G'\tau_{\beta+1}$ . Also,  $y \in Tx \rightarrow (\exists w \in M_{\beta}, y \approx w) \vee y \in \bigcup_{\gamma < \beta} K_{\gamma}$ . Since  $x \in M_{\beta}$ ,  $\exists z \in M_{\beta}, x \approx y$ .  $\therefore x \in M_{\beta+1}$ , i.e.,  $x \in K_{\beta} \cdot M_{\beta+1}$ ,  $\rightarrow$  induction hypothesis.

3.  $\delta \neq \beta \rightarrow K_{\delta} \subseteq K_{\beta}$ .

Proof:  $x \in K_{\delta} \rightarrow x \in R_{\tau_{\delta}} \subseteq R_{\tau_{\beta}}$ . By 2.,  $x \notin M_{\beta}$ .

4.  $v \approx w \rightarrow v \in W$  (where  $V$  is the universe of  $\text{NBG} + (\text{AI})$  and  $v \in V$ ).

Proof: If  $x \in v$ ,  $\exists$  least  $\lambda, x \in M_{\lambda} \cup K_{\lambda}$ . By the axiom of replacement these  $\beta$  form a set. Let  $\lambda$  be an upper bound of such  $\beta$ . Then  $x \in v \rightarrow x \in M_{\lambda} \cup K_{\lambda} \subseteq M_{\lambda} \cup K_{\lambda}$  by 1. and 3., i.e.,  $v \approx M_{\lambda} \cup K_{\lambda}$ . Let  $y \in v$ . If  $y \in M_{\lambda}$ ,  $\exists w \in M_{\beta}, y \approx w$ ; if  $y \in K_{\lambda}$ ,  $y \in \bigcup_{\gamma < \lambda} K_{\gamma}$ .  $\therefore y \in v \rightarrow (\exists w \in M_{\beta}, y \approx w) \vee y \in \bigcup_{\gamma < \lambda} K_{\gamma}$ . Let  $y \in \bigcup_{\gamma < \lambda} (U_{\gamma} \cup V)$ ,  $n > 0$ . Then  $\exists z \in v, y \in \bigcup_{m=1}^{n-1} U_z$ . Since  $y \in Tz$  and  $z \in M_{\lambda} \cup K_{\lambda}$ ,  $(\exists w \in M_{\beta}, y \approx w) \vee y \in \bigcup_{\gamma < \lambda} K_{\gamma}$ .  $\therefore y \in Tv \rightarrow (\exists w \in M_{\beta}, y \approx w) \vee y \in \bigcup_{\gamma < \lambda} K_{\gamma}$ . Since  $v \approx M_{\lambda} \cup K_{\lambda}$ ,  $v \in R_{\tau_{\lambda+1}} = G'\tau_{\lambda+1}$ .  $\therefore v \in G'\tau_{\lambda+1} + 1 \in \bigcup G''\tau_{\lambda+1} = R_{\tau_{\lambda+1}}$ .  $\therefore v \in R_{\tau_{\lambda+2}}$ . If  $\exists w \in M_{\beta}, v \approx w$  then  $v \in M_{\lambda+2}$ ; otherwise  $v \in K_{\lambda+2}$ .

5.  $V = W$ . (Hence  $w$  satisfies the general axioms.)

Proof: If  $V - W \neq \emptyset$ ,  $\exists y \in V - W, y \cdot (V - W) = \emptyset$ . Hence by 4.,  $y \in W$ ,  $\rightarrow$ .

6.  $x \approx y \in M_{\beta} \rightarrow x \in M_{\beta}$  (Subsets<sub>m</sub>).

Proof: By  $(\Phi_{\beta})$ ,  $x \in R_{\tau_{\beta}}$ .  $\exists z \in M_{\beta}, y \approx z$ .  $\therefore \exists w \in z, x \approx w$ , and  $w \in M_{\beta}$  by  $(\Phi_{\beta})$ . If  $z \in Tx$  then  $z \in Ty$ , so  $(\Phi_{\beta})$  holds for  $x$ .

7.  $x, y \in M_{\beta} \rightarrow x \cdot y \in M_{\beta}$ .

Proof: By  $(\Phi_{\beta})$ ,  $x \cdot y \in R_{\tau_{\beta}}$ .  $\exists u, v \in M_{\beta}, u \approx x, v \approx y$ . By  $(\Phi_{\beta})$ ,  $u \cdot v \in M_{\beta}$ . If  $z \in T(x \cdot y)$  then  $z \in Tx \cdot Ty$ , so  $(\Phi_{\beta})$  holds for  $x \cdot y$ .

8.  $x, y \in M_{\beta} \rightarrow \{x, y\} \in M_{\beta}$  (Pairing<sub>m</sub>).

Proof: By  $(\Phi_{\beta})$ ,  $\{x\}, \{y\} \in R_{\tau_{\beta}}$ . Obviously,  $(\Phi_{\beta})$  holds for  $\{x\}$  and  $\{y\}$ .

$\therefore \{x, y\} \in M_{\beta}$  by 7.

9.  $x \in M_{\beta} \rightarrow \bigcup_{\gamma < \beta} (x \cdot \bigcup M_{\gamma}) \in M_{\beta}$  (Union<sub>m</sub>).

proof: Since  $x \cdot U M_\gamma \in x$ ,  $x \cdot U M_\gamma \in R_{\tau_\beta}$  by  $(\Phi_\beta)$ .  $\therefore U (x \cdot U M_\gamma) \in R_{\tau_\beta}$  by  $(\Psi_\beta)$ . Now  $U (x \cdot U M_\gamma) \subseteq U \bar{z}$ . For each such  $z$ ,  $z \in U M_\gamma \rightarrow z \in \tau_\beta$ . Also,  $x \cdot U M_\gamma \in M_\beta$  by  $\circ$ , so  $x \cdot U M_\gamma \in \tau_\beta$ . Hence by regularity,  $U \bar{z} \in \tau_\beta$ . Since  $T(U (x \cdot U M_\gamma)) \subseteq T(x \cdot U M_\beta)$ ,  $(\Phi_\beta)$  for  $U (x \cdot U M_\beta)$  follows from  $(\Phi_\beta)$  for  $x \cdot U M_\beta$ . But the latter follows from 6.

10.  $x \in M_\beta \rightarrow Px \in M_\beta$  (Power $_\beta$ ).

Proof:  $Px \in R_{\tau_\beta}$  by  $(\Psi_\beta)$ .  $\exists y \in R_{\tau_\beta} \rightarrow x \cdot y$ . By  $(\Psi_\beta)$ ,  $Px \in R_{\tau_\beta}$ . Also,  $T(Px) = Px \in Tx$ .

$\therefore (\Phi_\beta)$  holds for  $Px$ .

11. Infinity $_\beta$ :  $\omega \in M_\beta$ .

12. Cardinality.

Proof: For any  $\beta$ ,  $x \in M_\beta \rightarrow \exists y \in M_\beta, x \cdot y$ , and  $Ty \in M_\beta$  by definition of  $M_\beta$ . Conversely, assume  $x \in W$  has the property  $\exists y \in U M_\beta, Ty \in U M_\beta$ , and  $x \cdot y$ . To show

$x \in U M_\beta$  it is necessary to show  $y \in M_\beta$ . Suppose not. Since  $y \in U M_\beta, y \notin K_\beta$ .  $\therefore y \notin M_\beta$ .  $\therefore \exists y_1 \in y$  with index  $\beta_1 > \tau_\beta$ . Since  $y_1 \in U M_\beta, y_1 \notin K_\beta$ . Since  $\beta_1 > \tau_\beta$ ,  $y_1 \notin M_\beta$ .  $\therefore y_1 \notin M_\beta$ .  $\therefore \exists y_2 \in y_1$  with index  $\beta_2 > \tau_\beta$ . Since  $y_2 \in U M_\beta, y_2 \notin K_\beta$ . Since  $\beta_2 > \tau_\beta$ ,  $y_2 \notin M_\beta$ .  $\therefore y_2 \notin M_\beta$ .  $\therefore \exists y_3 \in y_2$  with index  $\beta_3 > \tau_\beta$ . Iteration yields descending  $\epsilon$ -sequences  $\dots \in y_2 \in y_1 \in y$ ,  $\dots \beta_3 < \beta_2 < \beta_1, \dots$ .  $\therefore y \in M_\beta$ .  $\exists$  minimal  $\beta, x \in M_\beta \cup K_\beta$ . If  $x \in K_\beta$  then  $x \notin R_{\tau_\beta}$ . For each  $w \in x$ ,  $\exists$  minimal  $i < \tau_\beta, w \in G^i$ . Since  $x \cdot y$  for some  $y \in K_\beta$ , these  $i$  form a  $\gamma$ -sequence for some  $\gamma < \tau_\beta$ . Let  $n$  be the l.u.b. of these  $i$ . Then  $x \in G^{n+1}$ , so  $x \in G^{n+1}$ . Since  $x \notin R_{\tau_\beta}, n \geq \tau_\beta$ ,  $\rightarrow$  regularity of  $\tau_\beta$ .  $\therefore x \in M_\beta$ .

13.  $z \in M_\beta \cup K_\beta \rightarrow z \in M_\beta \cup K_\beta$ .

Proof:  $z \in M_\beta \cup K_\beta \rightarrow z \in R_{\tau_\beta}$ .  $\therefore y \in z \rightarrow y \in R_{\tau_\beta} \therefore y \in R_{\tau_\beta}$ . Also,  $Ty \in Tx$ . If  $\exists w \in M_\beta, w \cdot y$ , then  $y \in M_\beta$ ; otherwise  $y \in K_\beta$ .  $\therefore z \in M_\beta \cup K_\beta$ .

14.  $z \in M_\beta \cup K_\beta \rightarrow \{x: Tx \subseteq U M_\beta \cup z\} \subseteq \{x: Tx \subseteq U M_\beta \cup U K_\beta\}$ .

Proof: Assume  $Tx \subseteq U M_\beta \cup z$ . Let  $y \in Tx$ . If  $y \notin U M_\beta$ , then  $y \in z$ , so  $y \in K_\beta$ , by 13.

15.  $\{x: Tx \subseteq U M_\beta \cup U K_\beta\} = \{x: Tx \subseteq M_\beta \cup U K_\beta\}$ .

Proof: Assume  $Tx \subseteq U M_\beta \cup U K_\beta$ . Suppose  $Tx \not\subseteq M_\beta \cup U K_\beta$ . Then  $\exists$  minimal  $\beta, \beta > \beta$ ,

$Tx \cdot M_\beta = \emptyset$ . Let  $x_1 \in Tx \cdot M_\beta$ , with  $x_1 \notin M_\beta$  for any  $\gamma < \beta$ . If  $x_1 \in R_{\tau_\beta}$  then  $x_1 \in K_\beta$ ,

$\rightarrow$ . (Since  $x_1 \in Tx$ ,  $(\bar{Q}_A)$  holds for  $x_1$  by 12.)  $\therefore \exists x_2 \in x_1, x_2 \notin R_{\tau_A}$ . If  $x_2 \in K_{\nu}$  for some  $\nu \in A$  then  $x_2 \in R_{\tau_{\nu}} = G^{\tau_{\nu}} \rightarrow \therefore x_2 \in G^{\tau_{\nu}} + 1 \in R_{\tau_A} \rightarrow \therefore x_2 \in \bigcup_{\nu \in \theta_{\sim}} M_{\nu}$ .  $\therefore x_2 \in M_{\beta_1}$  for some minimal  $\beta_1$ . Since  $x_1 \in x_1$ ,  $\beta_1 \leq \beta_0$ . Since  $\beta_0$  is minimal,  $Tx \cdot M_{\beta_0} \neq \emptyset$ ,  $\beta_1 = \beta_0$ . If  $x_2 \in R_{\tau_A}$  then  $x_2 \in K_A \rightarrow \therefore \exists x_3 \in x_2, x_3 \notin R_{\tau_A}$ . By the same argument, with  $x_1$  and  $x_2$  replaced by  $x_2$  and  $x_3$ , respectively,  $x_3 \in M_{\beta_1}$ . Iteration of this argument yields an infinite descending  $\epsilon$ -sequence of elements of  $M_{\beta_0}$ ,  $\dots \rightarrow x_3 \in x_2 \in x_1 \rightarrow \dots$ .

16.  $M_A \cup \bigcup_{\nu \in A} K_{\nu}$  is transitive.

Proof: Let  $Y_A = M_A \cup \bigcup_{\nu \in A} K_{\nu}$ ,  $y \in x \in Y_A$ . If  $x \in M_A$ , since  $y \in Tx$ , either  $\exists w \in M_A, yzw$  or  $y \in \bigcup_{\nu \in A} K_{\nu}$ . In the first case  $y \in \bigcup_{\nu \in \theta_{\sim}} M_{\nu}$  by 12. Since  $x \in R_{\tau_A}$ ,  $y \in R_{\tau_A}$ . Since  $Ty \subseteq Tx$ ,  $y \in M_A$ .  $\therefore y \in Y_A$ . Suppose  $x \in K_{\beta}$  for some  $\beta \in A$ . Either  $\exists w \in M_A, yzw$  or  $y \in \bigcup_{\nu \in \beta} K_{\nu} \subseteq \bigcup_{\nu \in A} K_{\nu}$ . In the first case  $y \in \bigcup_{\nu \in \theta_{\sim}} M_{\nu}$  by 12. Since  $x \in R_{\tau_{\beta}}$ ,  $y \in R_{\tau_{\beta}} \subseteq R_{\tau_A}$ . Since  $Ty \subseteq Tx$ ,  $y \in M_A$ .  $\therefore y \in Y_A$ .

17.  $M_A, K_A \in W$ .

Proof:  $M_A \in R_{\tau_A}$ , and  $M_A \notin M_A$ . To show  $M_A \in K_A$ , it suffices to show  $TM_A \subseteq Y_A = M_A \cup \bigcup_{\nu \in A} K_{\nu}$ . Let  $x \in TM_A$ . If  $x \in M_A$  then  $x \in Y_A$ . If  $x \in \bigcup_{\nu \in A} M_{\nu}$ ,  $\exists y \in M_A, x \in y$ . By 16.,  $x \in y \in Y_A$ . If  $x \in \bigcup_{\nu \in \theta_{\sim}} M_{\nu}$ ,  $\exists y_1 \in \bigcup_{\nu \in \theta_{\sim}} M_{\nu}, x \in y_1$ , and  $\exists y_2 \in M_A, y_1 \in y_2$ . By 16.,  $y_1 \in y_2 \in Y_A$  and  $x \in y_1 \in Y_A$ . In general, if  $x \in \bigcup_{\nu \in \theta_{\sim}} M_{\nu}$  then  $x \in Y_A$  by  $n$  applications of 16.  $\therefore M_A \in W$ . By definition,  $K_A \subseteq PR_{\tau_A} = G^{\tau_A} + 1 \subseteq G^{\tau_{A+1}}$ , i.e.,  $K_A \in PR_{\tau_{A+1}}$ . Also,  $K_A \subseteq K_{\theta}$ , so  $K_A \subseteq K_{\beta} \in M_{\beta}$ .  $\therefore K_A \in \bigcup_{\nu \in \theta_{\sim}} M_{\nu}$ . Let  $x \in TK_A$ . If  $x \in K_A$  then  $x \in Y_{A+1}$ . If  $x \in \bigcup_{\nu \in \theta_{\sim}} K_{\nu}$  then  $\exists y \in K_{\beta}, x \in y$ . Since  $y \in K_{\beta} \subseteq Y_{\beta+1}$ ,  $x \in Y_{\beta+1}$  by 16. If  $x \in \bigcup_{\nu \in \theta_{\sim}} K_{\nu}$  then  $x \in y_1 \in y_2$  with  $y_1 \in \bigcup_{\nu \in \theta_{\sim}} K_{\nu}$  and  $y_2 \in K_{\beta} \subseteq Y_{\beta+1}$ .  $\therefore y_1 \in Y_{\beta+1}$  by 16., so  $x \in Y_{\beta+1}$  by 16. In general  $TK_A \subseteq Y_{A+1}$ .  $\therefore K_A \in K_{A+1}$ .

18.  $\{x: Tx \subseteq M_A \cup \bigcup_{\nu \in A} K_{\nu}\} \subseteq M_A \cup K_A$ .

Proof: Assume  $Tx \subseteq M_A \cup \bigcup_{\nu \in A} K_{\nu}$ . Let  $y \in x$ . If  $y \in M_A$  then  $y \in K_{\tau_A}$ . Otherwise  $y \in K_{\nu}$  for some  $\nu \in A$ , so  $y \in R_{\tau_{\nu}}$ .  $\therefore y \in G^{\tau_{\nu}} + 1 \in R_{\tau_A}$ .  $\therefore x \in R_{\tau_A}$ . Also,  $(\bar{Q}_A)$  holds for  $x$ . Hence if  $x \notin M_A$  then  $x \in PR_{\tau_A} = M_A$ , i.e.,  $x \in K_A$ .

19.  $\{x: Tx \subseteq \bigcup_{\nu \in \theta_{\sim}} M_{\nu} \cup z\} \in V$  (Universe  $\omega$ ).

Proof: By 14., 15., 18., 17.

$$20. M_{\alpha} = \{x \in \bigcup_{\gamma < \alpha} M_{\gamma} : \exists y \in \bigcup_{\gamma < \alpha} M_{\gamma} \cup \bigcup_{\gamma < \alpha} K_{\gamma}\}, K_{\alpha} = \{x \in W - \bigcup_{\gamma < \alpha} M_{\gamma} : \exists y \in \bigcup_{\gamma < \alpha} M_{\gamma} \cup \bigcup_{\gamma < \alpha} K_{\gamma}\}.$$

Proof: Let  $x \in M_{\alpha} \cup K_{\alpha}$ ,  $y \in Tx$ . If  $\exists w \in M_0, y \neq w$ , then  $y \in \bigcup_{\gamma < \alpha} M_{\gamma} \cup \bigcup_{\gamma < \alpha} K_{\gamma}$ . Conversely, assume  $Tx \in \bigcup_{\gamma < \alpha} M_{\gamma} \cup \bigcup_{\gamma < \alpha} K_{\gamma}$ . Then  $Tx \in M_{\alpha} \cup K_{\alpha}$ , by 15., so  $x \in M_{\alpha} \cup K_{\alpha}$  by 18.

9. Some immediate relations between K and other systems.

1. Let A be a formula of K which does not involve M (and thus does not involve any  $M_{\alpha}$  or  $K_{\alpha}$ ). Then A may also be regarded as a formula of NBG (although the variables of A as a formula of K and as a formula of NBG have different ranges). If  $\vdash_{NGB} A$  then  $\vdash_K A$ , for a proof in NBG is a proof in K. But the converse need not hold.

Proof: In K it can be shown that  $\exists x$  which determines a supercomplete model for NBG, viz.,  $x = M_0$ . In NBG this is not provable, for it would imply the existence of an inaccessible ordinal.

2. Axiom 7 of [1] does not hold for the system K, since  $K_{\alpha} \neq M_{\alpha}$ , follows from Axiom 7 of [1] (theorem 2.9 in [1]).

3. The theory K does not satisfy the axioms of the theory  $G_{\omega}$  of [4], for as observed in [1],  $V \neq V$  in  $G_{\omega}$ .

4. Suppose the model  $\Delta$  for NBG+(GCH) described in [2] is constructed in K. Extend  $\Delta$  to an interpretation of K by taking  $\mathfrak{M}_{\alpha}$  (the multiplicity of extended sets of the interpretation) to be  $\aleph \cdot L$ . Then  $\Delta$  is a model for K. (Hence the axiom of constructibility  $V=L$  and  $\Delta$ -GCH are consistent with K.)

Proof: The notation is as in [2] and the numbers in parentheses refer to statements in [2]. All notions, operations, classes and variables occurring in the special axioms are shown in [2] to be absolute, except for the power class operation; hence the subscript  $\Delta$  denoting relativization to  $\Delta$  is omitted except for the power class operation.

(1)  $\aleph \cdot L$  is a constructible multiplicity.

Proof: It is required to show  $F^{\aleph} \cdot \aleph \cdot L \leq L$  for all  $\aleph \in \text{On}(9.41)$ . But  $F^{\aleph} \leq L$  (9.51), so it is required to show  $\aleph \cdot F^{\aleph} \leq L$ . Assume for ordinals  $\alpha < \aleph$  and

Let  $\alpha = J_i \langle \beta, \gamma \rangle$ ,  $0 \leq i \leq 3$ . Then  $\alpha > \text{Max}\{\beta, \gamma\}$  if  $i \neq 0$  (9.25).

(a)  $i=0$ . Then  $F'_\alpha = F''_\alpha$  (9.35), so  $F'_\alpha \in M$ .

(b)  $i=1$ . Then  $F'_\alpha = \{F'_\beta, F'_\gamma\}$  (9.31), so  $F'_\alpha \in M$ .

(c)  $i=2$ . Then  $F'_\alpha = E \cdot F'_\beta$  (9.32). But  $E \in M$ , so  $M \cdot F'_\alpha = F'_\alpha$ .

(d)  $i=3$ . Then  $F'_\alpha = F'_\beta - F'_\gamma$  (9.33).  $\therefore M \cdot F'_\alpha = (M \cdot F'_\beta) - F'_\gamma \in L$  (9.6,  $i=3$ ).

(e)  $4 \leq i \leq \delta$ . Then  $F'_\alpha = F'_\beta \cdot J_i''(F'_\gamma)$  (9.34), so  $M \cdot F'_\alpha = (M \cdot F'_\beta) \cdot J_i''(F'_\gamma)$ . By induction hypothesis,  $E \in M$ ,  $M \cdot F'_\beta = F'_\beta$ , so  $M \cdot F'_\alpha = F'_\alpha$   $J_i \langle \gamma, \delta \rangle \in L$ .

(ii) Subsets $_M$ :  $x \in M \cdot L, Z \in \mathcal{L}(Z) \rightarrow x \cdot Z \in M \cdot L$ .

Proof: By Subsets $_M$ ,  $x \cdot Z \in M$ , and by definition of  $\mathcal{L}(Z)$ . (9.41),  $x \cdot Z \in L$ .

(iii) Union $_M$ :  $x \in M \cdot L \rightarrow \cup x \cdot M \cdot L \in M \cdot L$ .

Proof: By Union $_M$ ,  $\cup x \cdot M \in M$ , so  $\cup x \cdot M \cdot L \subseteq \cup x \cdot M \rightarrow \cup x \cdot M \cdot L \in M$  by Subsets $_M$ .

Since  $\mathcal{L}(M \cdot L)$  by (i),  $x \cdot M \cdot L \in L$  (9.41), so  $\cup x \cdot M \cdot L \in L$  since  $\Delta$  is a model for NBG.

(iv) Pairing :  $x, y \in M \cdot L \rightarrow \{x, y\} \in M \cdot L$ .

Proof: By Pairing $_M$ ,  $\{x, y\} \in M$ ; since  $\Delta$  is a model for NBG,  $\{x, y\} \in L$ .

(v) Infinity $_M$ :  $(\exists x)(x \in M \cdot L \wedge \{x\} \in M \cdot L \wedge x$  is infinite).

Proof: The construction of  $K$  from the extended von Neumann function shows  $\omega \in M \cdot L \wedge \omega \in M$ . By (11.42),  $\omega \in \text{On} \subseteq L$ .

(vi) Universes:  $z \in L \rightarrow \{y \in L : Ty \subseteq M \cdot L \cdot z\} \in L$ .

Proof: Let  $x = \{y \in L : Ty \subseteq M \cdot L \cdot z\}$ . By  $K$ ,  $x \in V$ , so  $E \in \text{On}$ ,  $\omega_\alpha > \text{Ord}'y$  for all  $y \in x$ .

Since  $\omega_\alpha \in R(J_\alpha)$  (\*9.27),  $x \subseteq F' \omega_\alpha$  (9.35).  $\therefore x \in F' \omega_{\alpha+1}$  (12.2).  $\therefore x \in L$  (9.51).

(vii) Cardinality:  $x \in L \rightarrow (x \in M \cdot L \leftrightarrow \exists y \in M \cdot L, Ty \subseteq M \cdot L \wedge x \approx y)$ .

Proof: If  $Ey \in M \cdot L, Ty \subseteq M \cdot L, x \approx y$ , then  $x \in M$  by  $K$ . Conversely, let  $x \in M \cdot L$ . By  $K$ ,  $Ez \in M$ ,  $x \approx z$ .  $E \in \text{On}$ ,  $z \in \alpha$ . But  $\text{On} \subseteq L$  (11.42), so that  $\alpha$  is the required value of  $y$ .

(viii) Power $_M$ :  $x \in M \cdot L \rightarrow P_x x \in M \cdot L$ .

Proof: Since  $x \in M$ ,  $Px \in M$  by Power $_M$ . But  $P_x x \subseteq Px$ , so  $P_x x \in M$  by Subsets $_M$ .

Since  $x \in L$ ,  $P_x x \in L$   $\therefore \Delta$  is a model for NBG.

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## AUTOBIOGRAPHICAL STATEMENT

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